

Chapter 1 – Historical and Economic Development of Computers

Computers are important to our society. Computers are everywhere. There – I have said the obvious. Many books on computer organization and architecture spend quite a few pages stating the obvious. This author does not want to waste paper on old news, so we proceed to our section on the history of computer development.

One might wonder why this author desires to spend time discussing the history of computers, since so much of what has preceded the current generation of computers is obsolete and largely irrelevant. In 1905, George Santayana (1863 – 1952) wrote that “Those who cannot remember the past are condemned to repeat it. ... This is the condition of children and barbarians, in whom instinct has learned nothing from experience.” [Santayana, *The Life of Reason*]. Those of us who do not know the past of computers, however, are scarcely likely to revert to construction with vacuum tubes and mercury delay lines. We do not study our history to avoid repeating it, but to understand the forces that drove the changes that are now seen as a part of that history.

Put another way, the issue is to understand the evolution of modern digital computers and the forces driving that evolution. Here is another way of viewing that question: A present-day computer (circa 2008) comprises a RISC CPU (to be defined later) with approximately 50 internal CPU registers, a pipelined execution unit, and one to four gigabytes of memory. None of this was unforeseeable in the 1940’s when the ENIAC was first designed and constructed, so why did they not build it that way. Indeed the advantages of such a unit would have been obvious to any of the ENIAC’s designers. During the course of reading this chapter we shall see why such a design has only recently become feasible. The history shows how engineers made decisions based on the resources available, and how they created new digital resources that drove the design of computers in new, but entirely expected, directions.

This chapter will present the history of computers in several ways. There will be the standard presentation of the computer generations, culminating in the question of whether we are in the fourth or fifth generation of computers. This and other presentations are linked closely with a history of the underlying technology – from mechanical relays and mercury delay lines, through vacuum tubes and core memory, to transistors and discrete components (which this author remembers without any nostalgia), finally to integrated circuits (LSI and VLSI) that have made the modern computer possible.

We shall then discuss the evolution of the computer by use, from single-user “bare iron” through single user computers with operating systems to multiple-user computer systems with batch and time-sharing facilities culminating in the variety of uses we see today.

Throughout this entire historical journey, the student is encouraged to remember this author’s favorite historical slogan for computers:

“The good old days of computing are today”

The Underlying Technologies

A standard presentation of the history of computing machines (now called computers) begins with a list of the generations of computers based on the technologies used to fabricate the CPU (Central Processing Unit) of the computer. These base technologies can be listed in four categories, corresponding somewhat to the generations that are to be defined later.

These technologies are:

- 1) Mechanical components, including electromechanical relays,
- 2) Vacuum tubes,
- 3) Discrete transistors, and
- 4) Integrated circuits, divided into several important classes.

To a lesser extent, we need to look at technologies used to fabricate computer memories, including magnetic cores and semiconductor memories.

Mechanical Components

All early computational devices, from the abacus to adding machines of the 1950's and early 1960's were hand-operated, with components that were entirely mechanical. These were made entirely of wood and/or metal and were quite slow in comparison to any modern electronic devices. The limitation of speed of such devices was due to the requirement to move a mass of material; thus incurring the necessity to overcome inertia. Those students who want to investigate inertia are requested either to study physics or to attempt simple experiments involving quick turns when running quickly (but not with scissors).

Analog and Digital Computers

With one very notable exception, most mechanical computers would be classified as analog devices. In addition to being much slower than modern electronic devices, these computers suffered from poor accuracy due to slippages in gears and pulleys. We support this observation with a brief definition and definition of the terms “analog” and “digital”.

Within the context of this discussion, we use the term “digital” to mean that the values are taken from a discrete set that is not continuous. Consider a clinical thermometer used to determine whether or not a small child has a fever. If it is an old-fashioned mercury thermometer, it is an analog device in that the position of the mercury column varies continuously. Another example would be the speedometer on most modern automobiles, it is a dial arrangement, where the indicator points to a number and moves continuously.

The author of these notes knows of one mechanical computer that can be considered to be digital. This is the abacus. As the student probably knows, the device records numbers by the position of beads on thin rods. While the beads are moved continuously, in general each bead must be in one of two positions in order for the abacus to be in a correct state.

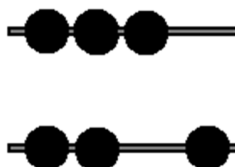


Figure: Abacus Bead Moves From One Position to The Other

Modern binary digital computers follow a similar strategy of having only two correct “positions”, thereby gaining accuracy. Such computers are based on electronic voltages, usually in the range 0.0 to 5.0 volts or 0.0 to 2.5 volts. Each circuit in the design has a rule for correcting slight imperfections in the voltage levels. In the first scheme, called TTL after the name of its technology (defined later), a voltage in the range 0.0 to 0.8 volts is corrected to 0.0 volts and a voltage in the range 2.8 to 5.0 volts is corrected to 5.0 volts. The TTL design reasonably assumes that voltages in the range 0.8 to 2.8 volts do not occur; should such a voltage actually occur, the behavior of the circuit will be unpredictable.

We have departed from our topic of mechanical computers, but we have a good discussion going so we shall continue it. In an analog electronic computer, a voltage of 3.0 would have a meaning different than a voltage of 3.1 or 2.9. Suppose that our intended voltage were 3.0 and that a noise signal of 0.1 volts has been added to the circuit. The level is now 3.1 volts and the meaning of the stored signal has changed; this is the analog accuracy problem. In the digital TTL world, such a signal would be corrected to 5.0 volts and hence its correct value.

A speedometer in an average car presents an excellent example of a mechanical analog computer, although it undoubtedly uses some electromagnetic components. The device has two displays of importance – the car speed and total distance driven. Suppose we start the automobile at time $T = 0$ and make the following definitions.

$S(T)$ the speed of the car at time T , and

$X(T)$ the distance the car has been driven up to time T .

As noted later, mechanical computers were often used to produce numerical solutions to differential and integral equations. Our automobile odometer provides a solution to either of the two equivalent equations.

$$\frac{dX(T)}{dT} = S(T)$$

$$X(T) = \int_0^T S(t) dt$$

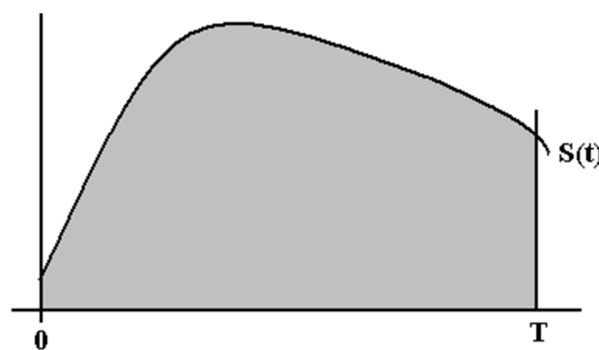


Figure: The Integral is the Area Under the Curve

The problem with such a mechanical computer is that slippage in the mechanism can cause inaccuracies. The speedometer on this author’s car is usually off by about 1.5%. This is OK for monitoring speeds in the range 55 to 70 mph, but not sufficient for scientific calculations.

Electronic Relays

Electronic relays are devices that use (often small) voltages to switch larger voltages. One example of such a power relay is the horn relay found in all modern automobiles. A small voltage line connects the horn switch on the steering wheel to the relay under the hood. It is that relay that activates the horn.

The following figure illustrates the operation of an electromechanical relay. The iron core acts as an electromagnet. When the core is activated, the pivoted iron armature is drawn towards the magnet, raising the lower contact in the relay until it touches the upper contact, thereby completing a circuit. Thus electromechanical relays are switches.

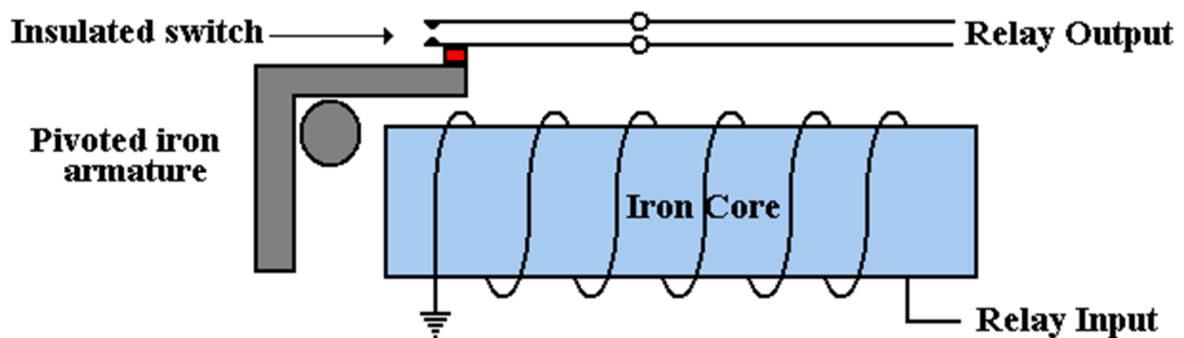
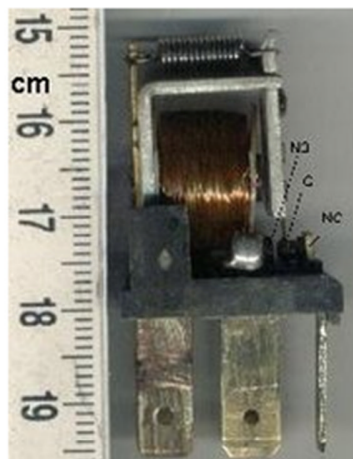
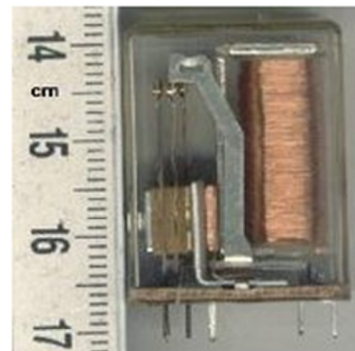


Figure: An Electromechanical Relay

In general, an electromechanical relay is a device that uses an electronic voltage to activate an electromagnet that will pull an electronic switch from one position to another, thus affecting the flow of another voltage; thereby turning the device “off” or “on”. Relays display the essential characteristic of a binary device – two distinct states. Below we see pictures of two recent-vintage electromechanical relays.



Relay for automotive use



Relay for use in electronics

Figure: Two Relays (Source <http://en.wikipedia.org/wiki/Relay>)

The primary difference between the two types of relays shown above is the amount of power being switched. The relays for use in general electronics tend to be smaller and encased in plastic housing for protection from the environment, as they do not have to dissipate a large amount of heat. Again, think of an electromechanical relay as an electronically operated switch, with two possible states: ON or OFF.

Power relays, such as the horn relay, function mainly to isolate the high currents associated with the switched apparatus from the device or human initiating the action. One common use is seen in electronic process control, in which the relays isolate the electronics that compute the action from the voltage swings found in the large machines being controlled.

In use for computers, relays are just switches that can be operated electronically. To understand their operation, the student should consider the following simple circuits.

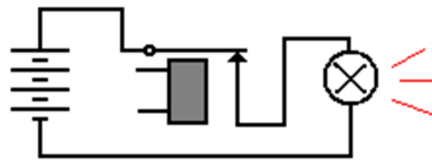


Figure: Relay Is Closed: Light Is Illuminated

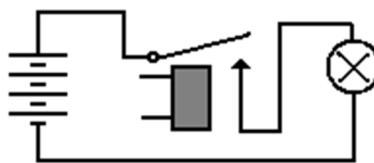


Figure: Relay Is Opened: Light Is Dark.

We may use these simple components to generate the basic Boolean functions, and from these the more complex functions used in a digital computer. The following relay circuit implements the Boolean AND function, which is TRUE if and only if both inputs are TRUE. Here, the light is illuminated if and only if both relays are closed.

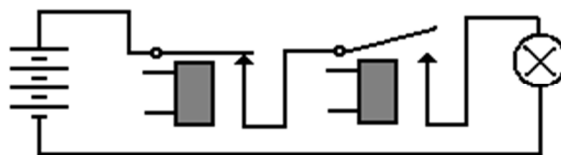


Figure: One Closed and One Open

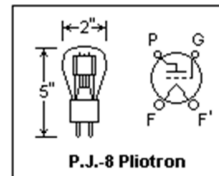
Computers based on electromagnetic relays played an important part in the early development of computers, but became quickly obsolete when designs using vacuum tubes (considered as purely electronic relays) were introduced in the late 1940's. These electronic tubes also had two states, but could be switched more quickly as there were no mechanical components to be moved.

Vacuum Tubes

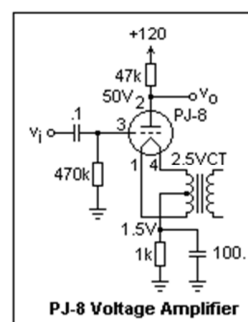
The next figure shows a picture of an early vacuum tube, along with its schematic and a circuit showing a typical use. The picture and schematics are taken from the web page of Dr. Elizabeth R. Tuttle of the University of Denver [R22].



Picture of Tube



Schematic of Tube



Use in a Circuit

The General Electric Plotron

This particular vacuum tube seems to have been manufactured in the 1920's. Later vacuum tubes (such as from the 1940's and 1950's) tended to be much smaller.

The vacuum tube shown above is a triode, which is a device with three major components, a cathode, a grid, and an anode. The cathode is the element that is the source of electrons in the tube. When it is heated either directly (as a filament in a light bulb) or indirectly by a separate filament, it will emit electrons, which either are reabsorbed by the cathode or travel to the anode. When the anode is held at a voltage more positive than the cathode, the electrons tend to flow from cathode to anode (and the current is said to flow from anode to cathode – a definition made in the early 19th century before electrons were understood). The third element in the device is a grid, which serves as an adjustable barrier to the electrons. When the grid is more positive, more electrons tend to leave the cathode and fly to the anode. When the grid is more negative, the electrons tend to stay at the cathode. By this mechanism, a small voltage change applied to the grid can cause a larger voltage change in the output of the triode; hence it is an amplifier. As a digital device, the grid in the tube either allows a large current flow or almost completely blocks it – “on” or “off”.

Here we should add a note related to analog audio devices, such as high-fidelity and stereo radios and record players. Note the reference to a small voltage change in the input of a tube causing a larger voltage change in the output; this is amplification, and tubes were once used in amplifiers for audio devices. The use of tubes as digital devices is just a special case, making use of the extremes of the range and avoiding the “linear range” of amplification.

We close this section with a picture taken from the IBM 701 computer, a computer from 1953 developed by the IBM Poughkeepsie Laboratory. This was the first IBM computer that relied on vacuum tubes as the basic technology.

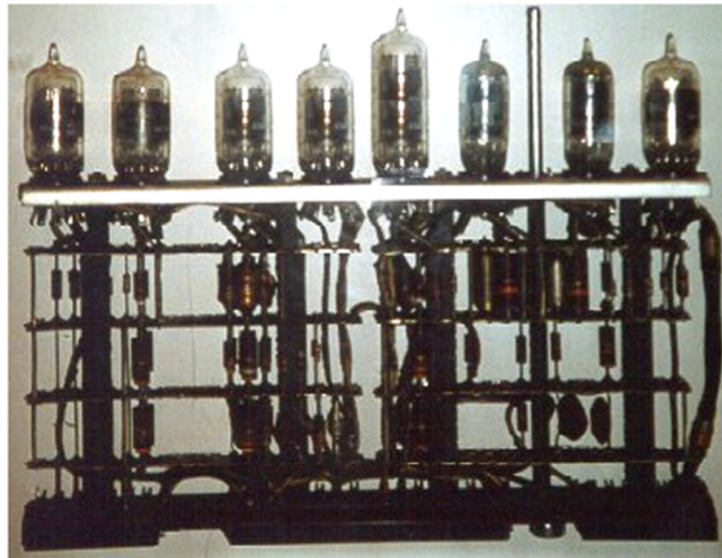


Figure: A Block of Tubes from the IBM 701

Source: [R23]

The reader will note that, though there is no scale on this drawing, the tubes appear smaller than the Pliotron, and are probably about one inch in height. The components that resemble small cylinders are resistors; those that resemble pancakes are capacitors.

Another feature of the figure above that is worth note is the fact that the vacuum tubes are grouped together in a single component. The use of such components probably simplified maintenance of the computer; just pull the component, replace it with a functioning copy, and repair the original at leisure.

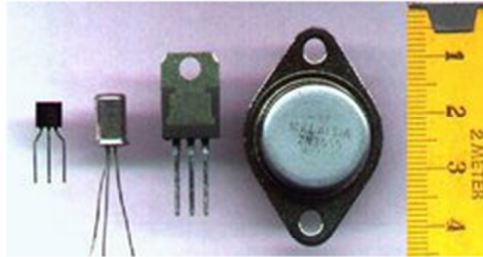
There are two major difficulties with computers fabricated from vacuum tubes, each of which arises from the difficulty of working with so many vacuum tubes. Suppose a computer that uses 20,000 vacuum tubes; this being only slightly larger than the actual ENIAC.

A vacuum tube requires a hot filament in order to function. In this it is similar to a modern light bulb that emits light due to a heated filament. Suppose that each of our tubes requires only five watts of electricity to keep its filament hot and the tube functioning. The total power requirement for the computer is then 100,000 watts or 100 kilowatts.

We also realize the problem of reliability of such a large collection of tubes. Suppose that each of the tubes has a probability of 99.999% of operating one hour. This can be written as a decimal number as $0.99999 = 1.0 - 10^{-5}$. The probability that all 20,000 tubes will be operational for more than an hour can be computed as $(1.0 - 10^{-5})^{20000}$, which can be approximated as $1.0 - 20000 \cdot 10^{-5} = 1.0 - 0.2 = 0.8$. There is an 80% chance the computer will function for one hour and 64% chance for two hours of operation. This is not good.

Discrete Transistors

Discrete transistors can be thought of as small triodes with the additional advantages of consuming less power and being less likely to wear out. The reader should note the transistor second from the left in the figure below. It at about one centimeter in size would have been an early replacement for the vacuum tube of about 5 inches (13 centimeters) in size shown on the previous page. One reason for the reduced power consumption is that there is no need for a heated filament to cause the emission of electrons.



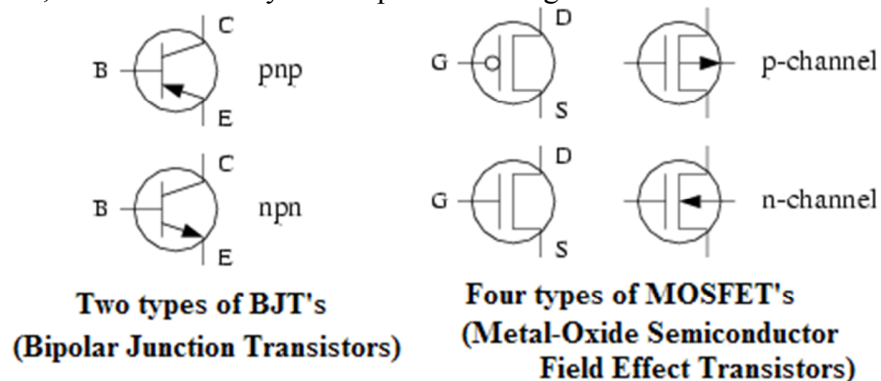
Source <http://en.wikipedia.org/wiki/Transistor>

When first introduced, transistors immediately presented an enormous (or should we say small – think size issues) advantage over the vacuum tube. However, there were cost disadvantages, as is noted in this history of the TX-0, a test computer designed around 1951. The TX-0 had 18-bit words and 3,600 transistors.

“Like the MTC [Memory Test Computer], the TX-0 was designed as a test device. It was designed to test transistor circuitry, to verify that a 256 X 256 (64-K word) core memory could be built ... and to serve as a prelude to the construction of a large-scale 36-bit computer. The transistor circuitry being tested featured the new Philco SBT100 surface barrier transistor, costing \$80, which greatly simplified transistor circuit design.” [R01, page 127]

That is a cost of \$288,000 just for the transistors. Imagine the cost of a 256 MB memory, since each byte of memory would require at least eight, and probably 16, transistors.

The figure below shows some of the symbols used to denote transistors in circuit diagrams. For this course, it is not necessary to interpret these diagrams.



Integrated Circuits

Integrated circuits are nothing more or less than a very efficient packaging of transistors and other circuit elements. The term “discrete transistors” implied that the circuit was built from transistors connected by wires, all of which could be easily seen and handled by humans.

Integrated circuits were introduced in response to a problem that occurred in designs that used discrete transistors. With the availability of small and effective transistors, engineers of the 1950's looked to build circuits of greater complexity and functionality. They discovered a number of problems inherent in the use of transistors.

- 1) The complexity of the wiring schemes. Interconnecting the transistors became the major stumbling block to the development of large useful circuits.
- 2) The time delays associated with the longer wires between the discrete transistors. A faster circuit must have shorter time delays.
- 3) The actual cost of wiring the circuits, either by hand or by computer aided design.

It was soon realized that assembly of a large computer from discrete components would be almost impossible, as fabricating such a large assembly of components without a single faulty connection or component would be extremely costly. This problem became known as the “tyranny of numbers”. Two teams independently sought solutions to this problem. One team at Texas Instruments was headed by Jack Kilby, an engineer with a background in ceramic-based silk screen circuit boards and transistor-based hearing aids. The other team was headed by research engineer Robert Noyce, a co-founder of Fairchild Semiconductor Corporation. In 1959, each team applied for a patent on essentially the same circuitry, Texas Instruments receiving U.S. patent #3,138,743 for miniaturized electronic circuits, and Fairchild receiving U.S. patent #2,981,877 for a silicon-based integrated circuit. After several years of legal battles, the two companies agreed to cross-license their technologies, and the rest is history (and a lot of profit).

The first integrated circuits were classified as SSI (Small-Scale Integration) in that they contained only a few tens of transistors. The first commercial uses of these circuits were in the Minuteman missile project and the Apollo program. It is generally conceded that the Apollo program motivated the technology, while the Minuteman program forced it into mass production, reducing the costs from about \$1,000 per circuit in 1960 to \$25 per circuit in 1963. Were such components fabricated today, they would cost less than a penny.

It is worth note that Jack Kilby was awarded the Nobel Prize in Physics for the year 2000 as a result of his contributions to the development of the integrated circuit. While his work was certainly key in the development of the modern desk-top computer, it was only one of a chain of events that lead to this development. Another key part was the development by a number of companies of photographic-based methods for production of large integrated circuits. Today, integrated circuits are produced by photo-lithography from master designs developed and printed using computer-assisted design tools of a complexity that could scarcely have been imagined in 1960.

Integrated circuits are classified into either four or five standard groups, according to the number of electronic components per chip. Here is one standard definition.

SSI	Small-scale integration	Up to 100 electronic components per chip Introduced in 1960.
MSI	Medium-scale integration	From 100 to 3,000 electronic components per chip Introduced in the late 1960's.
LSI	Large-scale integration	From 3,000 to 100,000 electronic electronic components per chip. Introduced about 1970.
VLSI	Very-large-scale integration	From 100,000 to 1,000,000 electronic components per chip Introduced in the 1980's.
ULSI	Ultra-large-scale integration	More than 1,000,000 electronic components per chip. This term is not standard and seems recently to have fallen out of use.

The figure below shows some typical integrated circuits. Actually, at least two of these figures show a Pentium microprocessor, which is a complete CPU of a computer that has been realized as an integrated circuit on a single chip. The reader will note that most of the chip area is devoted to the pins that are used to connect the chip to other circuitry.

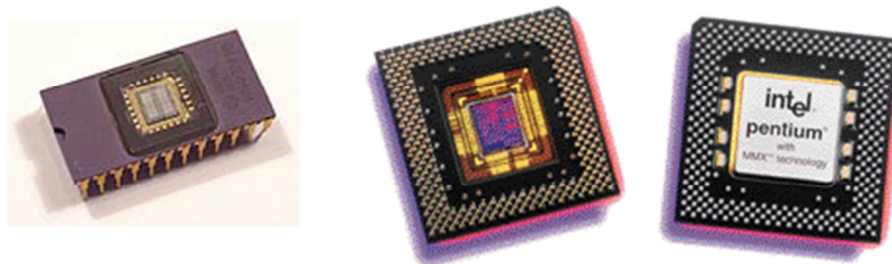


Figure: Some Late-Generation Integrated Circuits

Note that the chip on the left has 24 pins used to connect it to other circuitry, and that the pins are arranged around the perimeter of the circuit. This chip is in the design of a DIP (Dual Inline Pin) chip. The Pentium (™ of Intel Corporation) chip has far too many pins to be arranged around its periphery, so that they are arranged along the surface of one side.

Having now described the technologies that form the basis of our descriptions of the generations of computer evolution, we now proceed to a somewhat historic account of these generations, beginning with the “generation 0” that until recently was not recognized.

The Human Computer

The history of the computer can be said to reach back in history to the first human who used his or her ability with numbers to calculate a number. While the history of computing machines must be almost as ancient as that of computers, it is not the same. The reason is that the definition of the word “computer” has changed over time. Before entering into a discussion of computing machines, we mention the fact that the word “computer” used to refer to a human being. To support this assertion, we offer a few examples.

Here are two definitions from the first edition of the Oxford English Dictionary [R10, Volume II]. The reader should note that the definitions shown for each word are the only definitions shown in the 1933 dictionary.

“Calculator – One who calculates; a reckoner”.

“Computer – One who computes; a calculator, reckoner; specifically a person employed to make calculations in an observatory, in surveying, etc.”

This last usage of the word “computer” is seen in a history of the Royal Greenwich Observatory in London, England. One of the primary functions of this observatory during the period 1765 to 1925 was the creation of tables of astronomical data; such work “was done by hand by human computers”. [R17]

We begin our discussion of human computers with a trio from the 18th century, three French aristocrats (Alexis-Claude Clairaut, Joseph-Jerome de Lalande, and Nicole-Reine Étable de la Brière Lepaute) who produced numerical solutions to a set of differential equations and computed the day of return of a comet predicted by the astronomer Halley to return in 1758. At the time, Newton’s theory of gravity had been proposed and commonly accepted, but regarded as not sufficiently tested. It was realized that a sufficiently detailed calculation of the gravitational interaction of the comet with the sun and major planets would yield a date of return that could be checked against observations in an attempt to verify Newton’s theory. In 1757 the trio predicted a return date of April 1758; the actual date was March 13, 1758.



Figure: Alexis-Claude Clairaut, Joseph-Jerome de Lalande, and Nicole-Reine Étable de la Brière Lepaute

By the end of the 19th century, the work of being a computer became recognized as an accepted professional line of work. This was undertaken usually by young men and boys with not too much mathematical training, else they would not follow instructions. Most of the men doing the computational work quickly became bored and sloppy in their work. For this reason, and because they would work more cheaply, women soon formed the majority of computers. The next figure is a picture of the “Harvard Computers” in 1913. The man in the picture is the director of the Harvard College Observatory.



Figure: The Computers at the Harvard College Observatory, May 13, 1913.

We continue our discussion of human computers by noting three women, two of whom worked as computers during the Second World War. Many of the humans functioning as computers and calculators during this time were women for the simple reason that many of the able-bodied men were in the military.

The first computer we shall note is Dr. Gertrude Blanch (1897–1996), whose biography is published on the web [R18]. She received a B.S. in mathematics with a minor in physics from New York University in 1932 and a Ph.D. in algebraic geometry from Cornell University in 1935. From 1938 to 1942, was the technical director of the Mathematical Tables Project in New York. By 1941 it employed 450 human computers. In 1942 the project became part of the wartime Office for Scientific Research and Development, and Blanch became a mathematician for the New York office of the Applied Math panel of the OSRD. During the war, the project calculated ballistics tables for use by the U.S. Army in aiming artillery and navigation tables for the U.S. Navy. In the late 1940's, Dr. Blanch became involved with the group that would later become the Association for Computing Machinery.

We note here that mathematical tables were quite commonly used to determine the values of most functions until the introduction of hand-held multi-function calculators in the 1990's.

At this point, the reader is invited to note that the name of the group was not “Association for Computers” (which might have been an association for the humans) but “Association for Computing Machinery”, reflecting the still-current usage of the word “computer”.

The next person we shall note is a woman who became involved with the ENIAC project at the Moore School of Engineering at the University of Pennsylvania. Kathleen McNulty was born in County Donegal, Ireland and came to the United States in 1924. She graduated with a degree in mathematics from Chestnut Hill College in Pennsylvania, one of the three women to graduate with that degree in 1942. She then joined the Moore School of Engineering as one of about 75 women employed as computers. The job of these human computers was to calculate tables used by soldiers in the U.S. Army to aim artillery shells (One web site has called this “the first killer app”). McNulty described her life as a “computer” as follows.

“We did have desk calculators at that time, mechanical and driven with electric motors that could do simple arithmetic. You’d do a multiplication and when the answer appeared, you had to write it down to reenter it into the machine to do the next calculation. We were preparing a firing table for each gun, with maybe 1,800 simple trajectories. To hand-compute just one of these trajectories took 30 or 40 hours of sitting at a desk with paper and a calculator. As you can imagine, they were soon running out of young women to do the calculations.” [R19]

Dr. Alan Grier of the George Washington University has studied the transition from human to electronic computers, publishing his work in a paper “Human Computers and their Electronic Counterparts” [R20], David Alan Grier from George Washington University analyzed the transition from human to electronic computers. In this paper, Dr. Grier notes that very few historians of computing bother with studying the work done before 1940, or as he put it.

“In studying either the practice or the history of computing, few individuals venture into the landscape that existed before the pioneering machines of the 1940s. Arguably, there is good reason for making a clean break with the prior era, as these computing machines seem to have little in common with scientific computation as it was practiced prior to 1940. At that time, the term ‘computer’ referred to a job title, not to a machine. Large-scale computations were handled by offices of human computers.”

In this paper, Dr. Grier further notes that very few of the women hired as human computers elected to make the transition to becoming computer programmers, as we currently use the term. In fact, the only women who did make the transition were the dozen or so who had worked on the ENIAC and were thus used to large scale computing machines. According to Dr. Grier [R20].

“Most human computers worked only with an adding machine or a logarithm table or a slide rule. During the 1940s and 1950s, when electronic computers were becoming more common in scientific establishments, human computers tended to view themselves as more closely to the scientists for whom they worked, rather than the engineers who were building the new machines and hence [they] did not learn programming [as the engineers did].”

Classification of Computing Machines by Generations

Before giving this classification, the author believes that he should cite directly a few web references that seem to present excellent coverage of the history of computers.

- 1) http://en.wikipedia.org/wiki/History_of_computing_hardware
- 2) http://nobelprize.org/physics/educational/integrated_circuit/history/
- 3) <http://ei.cs.vt.edu/~history/index.html>
- 4) <http://www.yorku.ca/sasit/sts/sts3700b/syllabus.html>

We take the standard classification by generations from the book *Computer Structures: Principles and Examples* [R04], an early book on computer architecture.

1. The first generation (1945 – 1958) is that of vacuum tubes.
2. The second generation (1958 – 1966) is that of discrete transistors.
3. The third generation (1966 – 1972) is that of small-scale and medium-scale integrated circuits.
4. The fourth generation (1972 – 1978) is that of large-scale integrated circuits..
5. The fifth generation (1978 onward) is that of very-large-scale integrated circuits.

This classification scheme is well established and quite useful, but does have its drawbacks. Two that are worth mention are the following.

1. It ignores the introduction of the magnetic core memory. The first large-scale computer to use such memory was the MIT Whirlwind, completed in 1952. With this in mind, one might divide the first generation into two sub-generations: before and after 1952, with those before 1952 using very unreliable memory technologies.
2. The term “fifth generation” has yet to be defined in a uniformly accepted way. For many, we continue to live with fourth-generation computers and are even now looking forward to the next great development that will usher in the fifth generation.

The major problem with the above scheme is that it ignores all work before 1945. To quote an early architecture book, much is ignored by this ‘first generation’ label.

“It is a measure of American industry’s generally ahistorical view of things that the title of ‘first generation’ has been allowed to be attached to a collection of machines that were some generations removed from the beginning by any reasonable accounting. Mechanical and electromechanical computers existed prior to electronic ones. Furthermore, they were the functional equivalents of electronic computers and were realized to be such.” [R04, page 35]

Having noted the criticism, we follow other authors in surveying “generation 0”, that large collection of computing devices that appeared before the official first generation.

Mechanical Ancestors (Early generation 0)

The earliest computing device was the abacus; it is interesting to note that it is still in use. The earliest form of abacus dates to 450 BC in the western world and somewhat earlier in China. The *suan-pan* (“computing tray”), a device similar to the modern abacus, seems to have been in general use in China since the 12th century CE.

For those who like visual examples, we include a picture of a modern abacus.

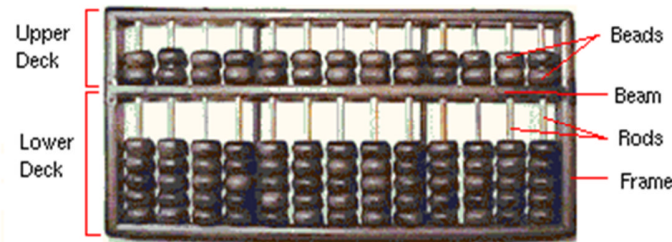


Figure: A Modern Abacus

In 1617, John Napier of Merchiston (Scotland) took the next step in the development of mechanical calculators when he published a description of his numbering rods, since known as ‘Napier’s bones’ for facilitating the multiplication of numbers. The set that belonged to Charles Babbage is preserved in the Science Museum at South Kensington, where there are other sets, some of which have come from the Napier family.”

The first real calculating machine, as the term is generally defined, was invented by the French philosopher Blaise Pascal in 1642. It was improved in 1671 by the German scientist Gottfried Wilhelm Leibniz, who developed the idea of a calculating machine which would perform multiplication by rapidly repeated addition. It was not until 1694 that his first complete machine was actually constructed. This machine is still preserved in the Royal Library at Hanover; close inspection shows that the mechanism for transmitting the carry part of the addition was never quite reliable. The first calculating machine successfully manufactured on a commercial scale was developed in 1820 by Charles Xavier Thomas of Colmar in Alsace.

A Diversion to Discuss Weaving – the Jacquard Loom

It might seem strange to divert from a discussion of computing machinery to a discussion of machinery for weaving cloth, but we shall soon see a connection. Again, the invention of the Jacquard loom was in response to a specific problem. In the 18th century, manual weaving from patterns was quite common, but time-consuming and plagued with errors. By 1750 it had become common to use punched cards to specify the pattern. At first, these instructions punched on cards were simply interpreted by the person (most likely a boy, as child labor was common in those days). Later modifications included provisions for mechanical reading of the pattern cards by the loom. The last step in this process was taken by Joseph Jacquard in 1801 when he produced a successful loom in which all power was supplied mechanically and control supplied by mechanical reading of the punched cards; thus becoming one of the first programmable machines controlled by punched cards.



The Jacquard Loom at the Deutsches Museum in Munich
Source: <http://65.107.211.206/technology/jacquard.html>

We divert from this diversion to comment on the first recorded incident of industrial sabotage. One legend concerning Jacquard states that during a public exhibition of his loom, a group of silk workers cast shoes, called “sabots”, into the mechanism in an attempt to destroy it. What is known is that the word “sabotage” acquired a new meaning, with this meaning having become common by 1831, when the silk workers revolted. Prior to about 1830, the most common meaning of the word “sabotage” was “to use the wooden shoes [sabots] to make loud noises”.

Charles Babbage and His Mechanical Computers

Charles Babbage (1792 – 1871) designed two computing machines, the difference engine and the analytical engine. The analytical engine is considered by many to be the first general-purpose computer. Unfortunately, it was not constructed in Babbage’s lifetime, possibly because the device made unreasonable demands on the milling machines of the time and possibly because Babbage irritated his backers and lost financial support.

Babbage’s first project, called the difference engine, was begun in 1823 as a solution to the problem of computing navigational tables. At the time, most computation was done by a large group of clerks (preferably not mathematically trained, as the more ignorant clerks were more likely to follow directions) who worked under the direction of a mathematician. Errors were common both in the generation of data for these tables and in the transcription of those data onto the steel plates used to engrave the final product.

The **difference engine** was designed to calculate the entries of a table automatically, using the mathematical technique of finite differences (hence the name) and transfer them via steel punches directly to an engraver's plate, thus avoiding the copying errors. The project was begun in 1823, but abandoned before completion almost 20 years later. The British government abandoned the project after having spent £17,000 (about \$1,800,000 in current money – see the historical currency converter [R24]) on it and concluding that its only use was to calculate the money spent on it.

The above historical comment should serve as a warning to those of us who program computers. We are “cost centers”, spending the money that other people, the “profit centers”, generate. We must keep these people happy or risk losing a job.



Figure: Modern Recreation of Babbage's Analytical Engine



Figure: Scheutz's Differential Engine (1853) – A Copy of Babbage's Machine

The reader will note that each of Babbage's machine and Scheutz's machine are hand cranked. In this they are more primitive than Jacquard's loom, which was steam powered.

Babbage's analytical engine was designed as a general-purpose computer. The design, as described in an 1842 report (<http://www.fourmilab.ch/babbage/sketch.html>) by L. F. Menebrea, calls for a five-component machine comprising the following.

- 1) *The Store* A memory fabricated from a set of counter wheels.
This was to hold 1,000 50-digit numbers.
- 2) *The Mill* Today, we would call this the Arithmetic-Logic Unit
- 3) *Operation Cards* Selected one of four operations: addition, subtraction, multiplication, or division.
- 4) *Variable Cards* Selected the memory location to be used by the operation.
- 5) *Output* Either a printer or a punch.

This design represented a significant advance in the state of the art in automatic computing for a number of reasons.

- 1) The device was general-purpose and programmable,
- 2) The provision for automatic sequence control,
- 3) The provision for testing the sign of a result and using that information in the sequencing decisions, and
- 4) The provision for continuous execution of any desired instruction.

To complete our discussion of Babbage's analytical engine, we include a program written to solve two simultaneous linear equations.

$$\begin{aligned} m \bullet x + n \bullet y &= d \\ m' \bullet x + n' \bullet y &= d' \end{aligned}$$

As noted in the report by Menebrea, we have $x = (d \bullet n' - d' \bullet n) / (m \bullet n' - m' \bullet n)$. The program to solve this problem, shown below, seems to be in assembly language.

Number of the operations	Operation-cards	Cards of the variables		Progress of the operations
	Symbols indicating the nature of the operations	Columns on which operations are to be performed	Columns which receive results of operations	
1	×	$V_2 \times V_4 =$	$V_8 \dots\dots$	$= dn'$
2	×	$V_5 \times V_1 =$	$V_9 \dots\dots$	$= d'n$
3	×	$V_4 \times V_0 =$	$V_{10} \dots\dots$	$= n'm$
4	×	$V_1 \times V_3 =$	$V_{11} \dots\dots$	$= nm'$
5	—	$V_8 - V_9 =$	$V_{12} \dots\dots$	$= dn' - d'n$
6	—	$V_{10} - V_{11} =$	$V_{13} \dots\dots$	$= n'm - nm'$
7	÷	$\frac{V_{12}}{V_{13}} =$	$V_{14} \dots\dots$	$= x = \frac{dn' - d'n}{n'm - nm'}$

Source: See R25

Augusta Ada Lovelace

No discussion of Babbage's analytical engine would be complete without a discussion of Ada Byron, Lady Lovelace. We have a great description from a web article written by Dr. Betty Toole [R26]. Augusta Ada Byron was born on December 10, 1815, the daughter of the famous poet Lord Byron. Five weeks after Ada's birth, Lady Byron asked for a separation from the poet and was awarded sole custody of her daughter.

When she was 17, Ada was introduced to Mary Somerville, who had translated the works of the French mathematician LaPlace into English. Mary encouraged Ada in her mathematical studies. In November 1834, Ada was introduced to Charles Babbage, at a dinner party hosted by Mary Somerville, at which Babbage discussed his analytical engine.

Babbage continued to work on plans for his analytical engine. In 1841, Babbage reported on progress at a seminar in Turin Italy. An Italian, Menebrea, attended this seminar and wrote a summary of Babbage's remarks. In 1843, Ada translated this article from the original French and showed the results to Babbage, who approved and suggested that she add her own notes. The resulting report was three times as long as Menebrea's original.

Ada suggested to Babbage a plan for how the analytical engine might calculate Bernoulli numbers. This plan is now regarded as the first program written for a general purpose computer. For this reason, Augusta Ada Lovelace is regarded as the world's first computer programmer. The programming language Ada was named in her honor in 1979.

Bush's Differential Analyzer

Some of the problems associated with early mechanical computers are illustrated by this picture of the Bush Differential Analyzer, built by Vannevar Bush of MIT in 1931. It was quite large and all mechanical, very heavy, and had poor accuracy (about 1%). [R21]

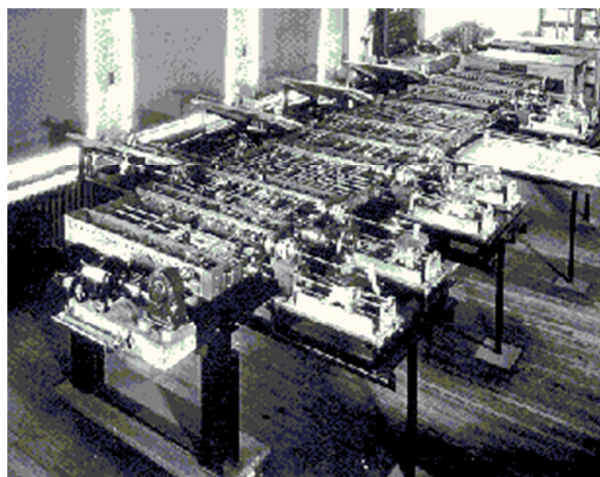


Figure: The Bush Differential Analyzer

For an idea of scale in this picture, the reader should note that the mechanism is placed on a number of tables (or lab benches), each of which would be waist high.

For the most part, computers with mechanical components were analog devices, set up to solve integral and differential equations. These were based on an elementary device, called a planimeter that could be used to produce the area under a curve. The big drawback to such a device was that it relied purely on friction to produce its results; the inevitable slippage being responsible for the limited accuracy of the machines.

A machine such as Bush's differential analyzer had actually been proposed by Lord Kelvin in the nineteenth century, but it was not until the 1930's that machining techniques were up to producing parts with the tolerances required to produce the necessary accuracy.

The differential analyzer was originally designed in order to solve differential equations associated with power networks but, as was the case with other computers, it was pressed into service for calculation of trajectories of artillery shells. There were five different copies of the differential analyzer produced; one was used at the Moore School of Engineering where it must have been contemporaneous with the ENIAC. The advent of electronic digital computers quickly made the mechanical analog computers obsolete, and all were broken up and mostly sold for scrap in the 1950's. Only a few parts have been saved from the scrap heap and are now preserved in museums. There is no longer a fully working mechanical differential analyzer from that period.

To finish this section on mechanical calculators, we show a picture of a four-function mechanical calculator produced by the Monroe Corporation about 1953. Note the hand cranks and a key used to reposition the upper part, called the "carriage".

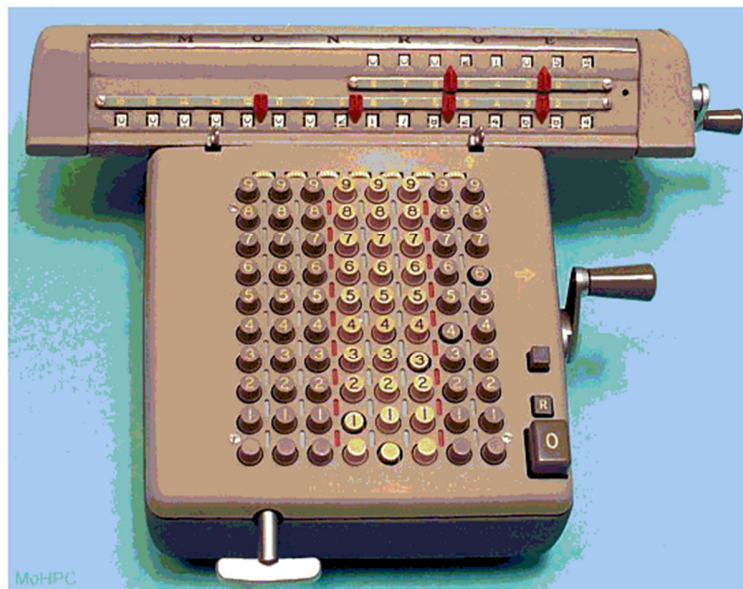


Figure: A Monroe Four-Function Calculator (circa 1953)

The purely mechanical calculators survived the introduction of the electronic computer by a number of years, mostly due to economic issues – the mechanicals were much cheaper, and they did the job.

Electromechanical Ancestors (Late generation 0 – mid 1930's to 1952)

The next step in the development of automatic computers occurred in 1937, when George Stibitz, a mathematician working for Bell Telephone Laboratories designed a binary full adder based on electromechanical relays. Dr. Stibitz's first design, developed in November 1937, was called the "Model K" as it was developed in his kitchen at home.

In late 1938, Bell Labs launched a research project with Dr. Stibitz as its director. The first computer was the Complex Number Calculator, completed in January 1940. This computer was designed to produce numerical solutions to the equations (involving complex numbers) that arose in the analysis of transmission circuits.

Stibitz worked in conjunction with a Bell Labs team, headed by Samuel Williams, to develop first a machine that could calculate all four basic arithmetic functions, and later to develop the Bell Labs Model Relay Computer, now recognized as the worlds first electronic digital computer. [R27]

Bell labs constructed six relay-based computers from 1938 through 1947. The characteristics of some of them are shown in the table below. Note "Model 11" for "Model II".

	Model 11	Model III	Model IV	Model V (two copies)
Date completed	7-1943	6-1944	3-1945	12-1946, 8-1947
Date dismantled	1961	1958	1961	1958
Place installed	Wash., D.C.	Ft. Bliss, Texas	Wash., D.C.	Langley, Va. Aberdeen, Md.
Also known as	Relay Interpolator	Ballistic Computer	Error Detector Mark 22	
No. relays	440	1,400	1,425	9,000+
Word length	2 to 5 fixed decimal	1 to 6 fixed pt.	1 to 6 fixed pt.	1 to 7 floating pt.
Memory cap.	7 numbers	10	10	30 total
Mult. speed	4 sec.	1 sec.	1 sec.	0.8 sec.
Mult. method	repeated add.	table look-up	table look-up	rep. add.
Cost	.	\$65,000	\$65,000	\$500,000
Size	2 panels	5 panels	5 panels	27 panels 10 tons

Table: The Model II – Model V Relay Computers [R09]

The Model V was a powerful machine, in many ways more reliable than and just as powerful as the faster electronic computers of its time. The relay computers became obsolete because the Model V represented the basic limit of the relay technology, whereas the electronic computers could be modified considerably.

The IBM Series 400 Accounting Machines

We now examine another branch of electromechanical computing devices that existed before 1940 and influenced the development of the early digital computers. This is a family of accounting machines produced by the International Business Machines Corporation (IBM), beginning with the IBM 405 Alphabetical Accounting Machine, introduced in 1934.

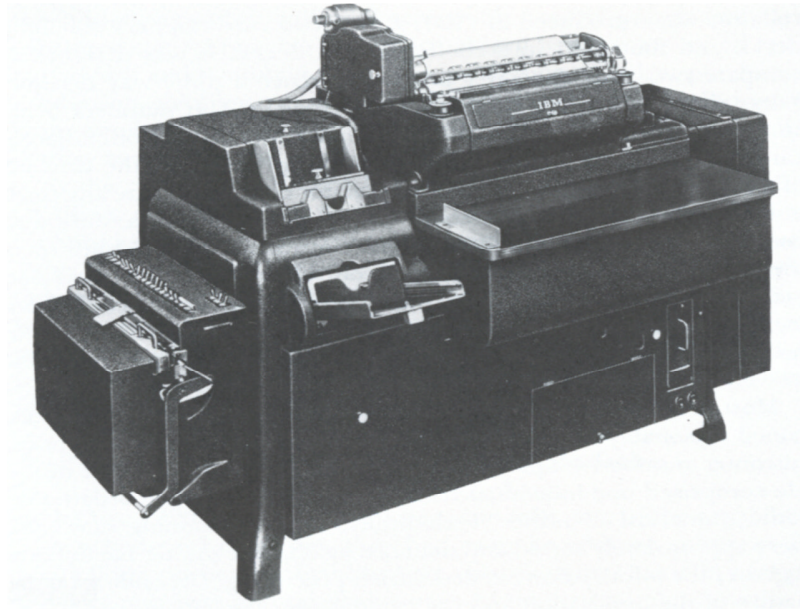


Figure: IBM Type 405 [R55]

According to the reference [R55], the IBM 405 was IBM's high-end tabulator offering (and the first one to be called an Accounting Machine). The 405 was programmed by a removable plugboard with over 1600 functionally significant "hubs", with access to up to 16 accumulators, the machine could tabulate at a rate of 150 cards per minute, or tabulate and print at 80 cards per minute. The print unit contained 88 type bars, the leftmost 43 for alphanumeric characters and the other 45 for digits only. The 405 was IBM's flagship product until after World War II (in which the 405 was used not only as a tabulator but also as the I/O device for classified relay calculators built by IBM for the US Army). In 1948, the Type 405 was replaced by an upgraded model, called the Type 402.

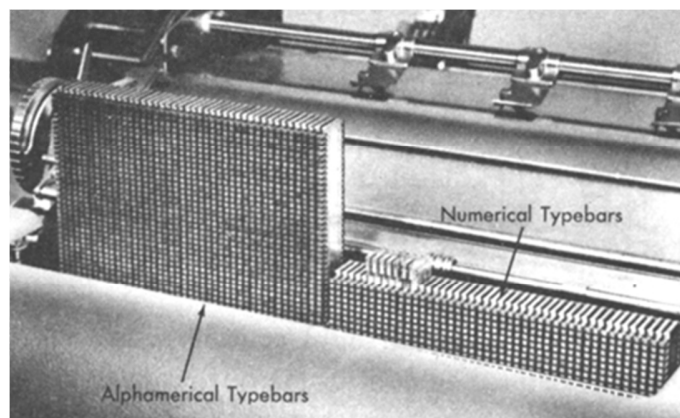


Figure: The type bars of an IBM 402 Accounting Machine [R55]

As noted above, the IBM 405 and its successor, the IBM 402, were programmed by wiring plugboards. In order to facilitate the use of “standard programs”, the plugboard included a removable wiring unit upon which a standard wiring could be set up and saved.

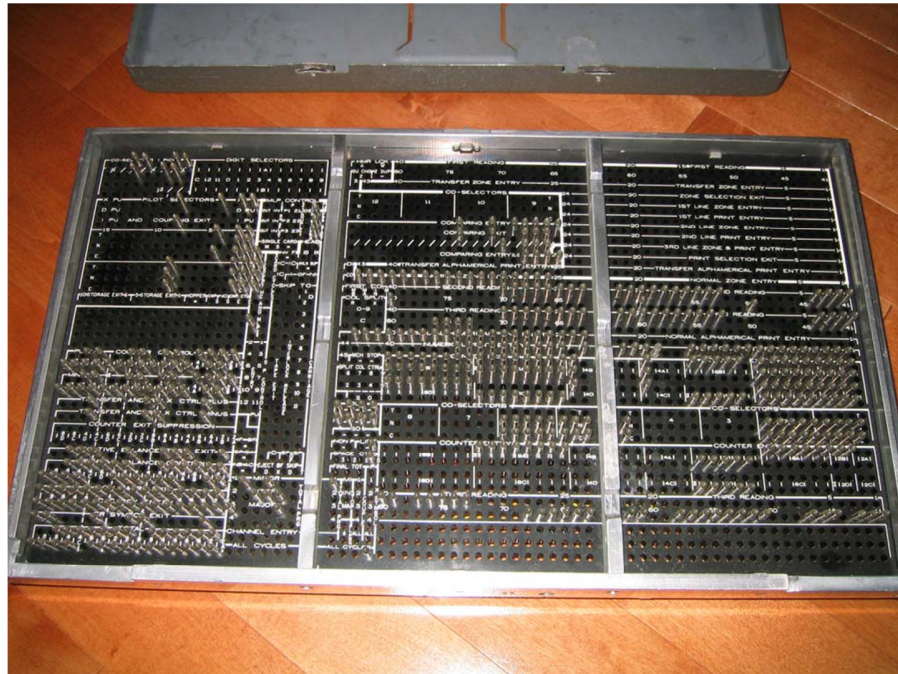


Figure: The Plugboard Receptacle

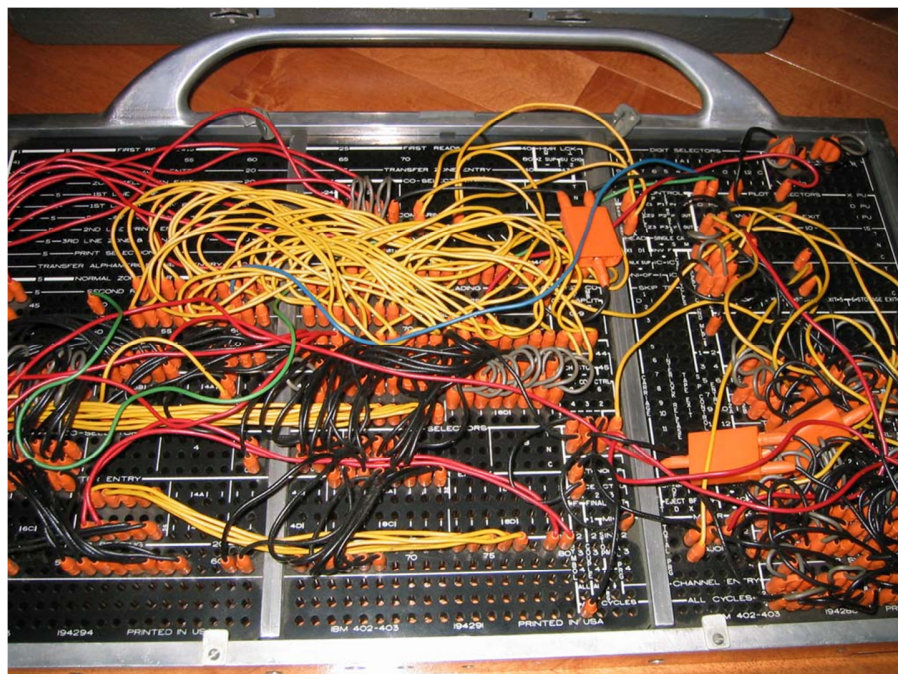


Figure: A Wired Plugboard

The Harvard Computers: Mark I, Mark II, Mark III, and Mark IV

The Harvard Mark I, also known as the IBM Automatic Sequence Controlled Calculator (ASCC) was the largest electromechanical calculator ever built. It has a length of 51 feet, a height of eight feet, and weighed nearly five tons. Here is a picture of the computer, from the IBM archives (http://www-1.ibm.com/ibm/history/exhibits/markI/markI_intro.html).



Figure: The Harvard Mark I

This computer was conceived in the 1930's by Howard H. Aiken, then a graduate student in theoretical physics at Harvard University. In 1939 IBM President Thomas J. Watson agreed to back the project, and it was presented to Harvard University on August 7, 1944; the delay being due to wartime demands associated with World War II.

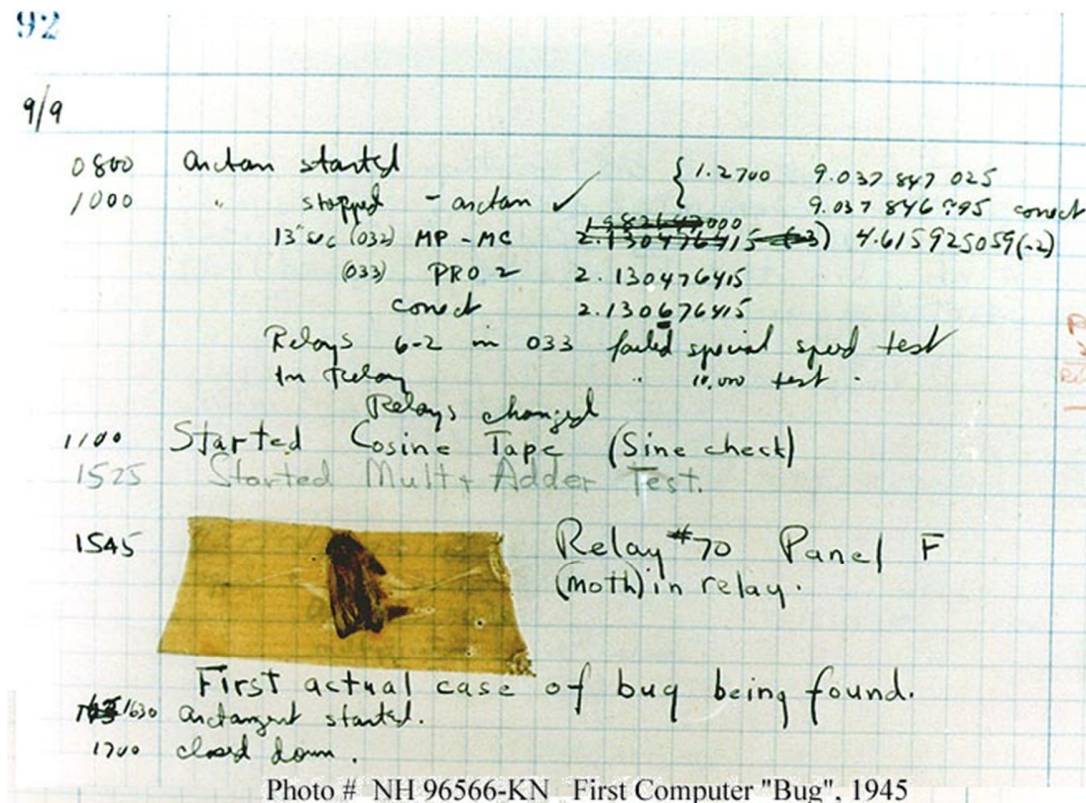
The Mark I operated at Harvard for 15 years, after which the machine was broken up and parts sent to the Smithsonian museum, a museum at Harvard, and IBM's collection. It was the first of a sequence of four electromechanical computers that lead up to the ENIAC.

The Mark II was begun in 1942 and completed in 1947. The Mark III, completed in 1949, was the first of the series to use an internally stored program and indirect addressing. The Mark IV, last of this series, was completed in 1952. It had a magnetic core memory used to store 200 registers and seems to have used some vacuum tubes as well.

IBM describes the Mark I, using its name "ASCC" (Automatic Sequence Controlled Calculator) in its archive web site, as

"consisting of 78 adding machines and calculators linked together, the ASCC had 765,000 parts, 3,300 relays, over 500 miles of wire and more than 175,000 connections. The Mark I was a parallel synchronous calculator that could perform table lookup and the four fundamental arithmetic operations, in any specified sequence, on numbers up to 23 decimal digits in length. It had 60 switch registers for constants, 72 storage counters for intermediate results, a central multiplying-dividing unit, functional counters for computing transcendental functions, and three interpolators for reading function punched into perforated tape."

The Harvard Mark II, the second in the sequence of electromechanical computers, was also built by Howard Aiken with support from IBM. In 1945 Grace Hopper was testing the Harvard Mark II when she made a discovery that has become historic – the "first computer bug", an occurrence reported in her research notebook, a copy of which is just below.



Source: Naval Surface Warfare Center Computer Museum [R28]

The caption associated with the moth picture on the web site is as follows.

“Moth found trapped between points at Relay # 70, Panel F, of the Mark II Aiken Relay Calculator while it was being tested at Harvard University, 9 September 1945. The operators affixed the moth to the computer log, with the entry: "First actual case of bug being found". They put out the word that they had ‘debugged’ the machine, thus introducing the term ‘debugging a computer program’.”

“In 1988, the log, with the moth still taped by the entry, was in the Naval Surface Warfare Center Computer Museum at Dahlgren, Virginia.”

We should note that this event is not the first occurrence of the use of the word “bug” to reference a problem or error. The term was current in the days of Thomas A. Edison, the great American inventor. One meaning of the word “bug” as defined in the second edition of the Oxford English Dictionary is “A defect or fault in a machine, plan, or the like, *origin U.S.*” What is of interest to us is when this meaning came into use. The OED provides a quotation from the March 11, 1889 edition of the Pall Mall Gazette, a publication very popular in England at the time.

“Mr. Edison, I was informed, had been up the two previous nights discovering ‘a bug’ in his phonograph – an expression for solving a difficulty and implying that some imaginary insect has secreted itself inside and is causing all the trouble.”

The Colossus

The Colossus was a computer built in Great Britain to attack the German diplomatic codes. We note here the common misconception that the Colossus was built to attack the German Naval codes generated using the coding machine known as Enigma. This is not so. The machines used to attack the ciphers generated by the Enigma machines were known as “Bombes” – large mechanical computers that used the strategy known as “brute force” to crack the Naval ciphers. These were called “Bombes” because the name resembled the Polish name for the machines. The English, in their attempt to disguise the Polish origin of the machines, encouraged the rumor that the machines were so called because they made a ticking noise not unlike the time-delay bombs dropped by the Germans on London.

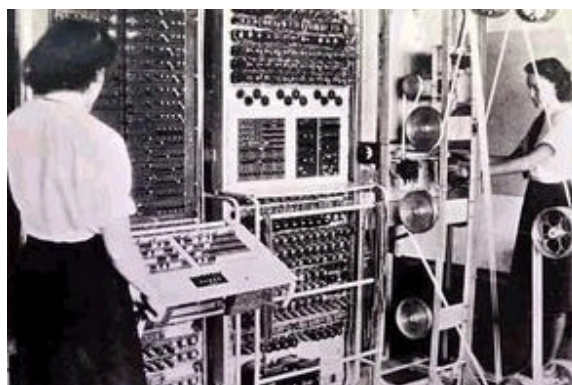


Figure: The Colossus

The reader should note the paper tape reels at the right of the picture.

Konrad Zuse (1910 – 1995) and the Z Series of Computers

Although he himself claimed not to have been familiar with the work of Charles Babbage, Konrad Zuse can be considered to have picked up where Babbage left off. Zuse used electromechanical relays instead of the mechanical gears that Babbage had chosen. Zuse built his first computer, the Z1, in his parent's Berlin living room; it was destroyed by allied bombing during the war and only later reconstructed (see below).

After the war, Zuse was employed in the United States and continued to develop computers. Some of his primary achievements are listed below.

Z1 (1938), a mechanical programmable digital computer. Although mechanical problems made its operation erratic, it was rebuilt by Zuse himself after the war.

Z2 (1940), an electro-mechanical computer.

Z3 (1941), this machine uses program control.

Plankalkül (1945/46), the first programming language, implemented on the Z3.

Z22 (1958), the last machine developed by Zuse. It was one of the first to be designed with transistors

The next figure shows the rebuilt Z1 in the Deutsche Technik Museum in Berlin. The picture was taken in 1989; from other sources this author has verified that the man standing by the display is Konrad Zuse himself.

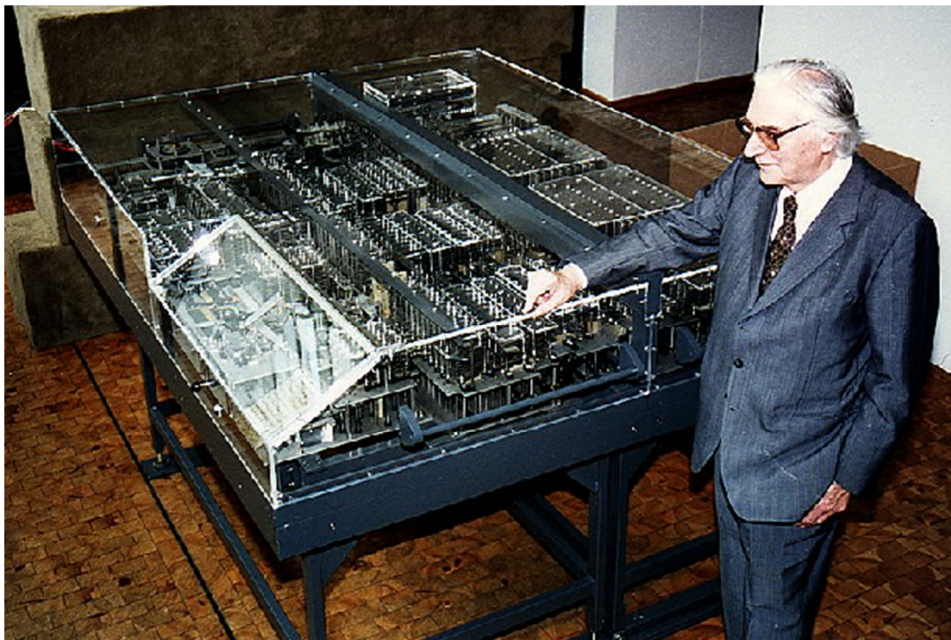


Figure: Konrad Zuse and the Reconstructed Z1
Source: R29

We use some of Konrad Zuse's own memoirs to make the transition to our discussion of the "first generation" of computers – those based on vacuum tube technology. Zuse begins with a discussion of Helmut Schreyer, a friend of his who first had the idea.

"Helmut was a high-frequency engineer, and on completing his studies (around 1936) ... suddenly had the bright idea of using vacuum tubes. At first I thought it was one of his student pranks - he was always full of fun and given to fooling around. But after thinking about it we decided that his idea was definitely worth a try. Thanks to switching algebra, we had already married together mechanics and electro-magnetics – two basically different types of technology. Why, then, not with tubes? They could switch a million times faster than elements burdened with mechanical and inductive inertia."

"The possibilities were staggering. But first basic circuits for the major logical operations such as conjunction, disjunction and negation had to be discovered. Tubes could not simply be connected in line like relay contacts. We agreed that Helmut should develop the circuits for these elementary operations first, while I dealt with the logical part of the circuitry. Our aim was to set up elementary circuits so that relay technology could be transferred to the tube system on a one-to-one basis. This meant the tube machine would not have to be redesigned from scratch. Schreyer solved this problem fairly quickly."

"This left the way open for further development. We cautiously told some friends about the possibilities. The reaction was anything from extremely skeptical to spontaneously enthusiastic. Interestingly enough, most criticism came from Schreyer's colleagues, who worked with tubes virtually all the time. They were doubtful that an apparatus with 2,000 tubes would work reliably. This critical attitude was the result of their own experience with large transmitters which contained several hundred tubes."

"After the War was finally over, news of the University of Pennsylvania ENIAC machine went all round the world – '18,000 tubes!'. We could only shake our heads. What on earth were all the tubes for?" [R30]

Vacuum Tube Computers (Generation 1 – from 1945 to 1958)

It is now generally conceded that the first purely electronic binary digital computer was designed by John Vincent Atanasoff (October 4, 1903 – June 15, 1995) at Iowa State University, with the assistance of Clifford Berry (April 19, 1918 – October 30, 1963). The computer, called the ABC for Atanasoff-Berry-Computer, used binary circuits built with vacuum tubes and a capacitive drum memory. Atanasoff and Berry worked on the computer from 1937 to 1942, when the demands of the war caused its development to stop. The first prototype of the computer was demonstrated to the university community in December 1939.

The relation of the ABC to the ENIAC (once claimed to be the first electronic digital computer) has been the source of much discussion. In fact, the claim of precedence of the ABC over the ENIAC was not settled until taken to court quite a bit later.

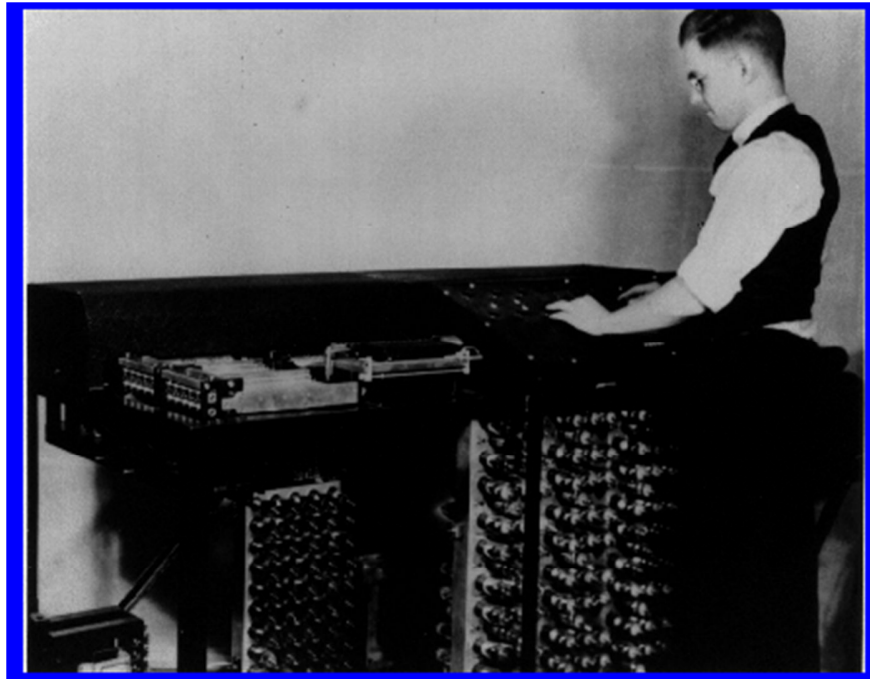


Figure: Clifford Berry standing in front of the ABC

Source: R31

Atanasoff [R 32] decided in building the ABC that:

- “1) He would use electricity and electronics as the medium for the computer.
- 2) In spite of custom, he would use base-two numbers (the binary system) for his computer.
- 3) He would use condensers (now called “capacitors” – see Review of Basic Electronics in these notes) for memory and would use a regenerative or "jogging" process to avoid lapses that might be caused by leakage of power.
- 4) He would compute by direct logical action and not by enumeration as used in analog calculating devices.”

According to present-day supporters of Dr. Atanasoff, two problems prevented him from claiming the credit for creation of the first general-purpose electronic digital computer: the slow process of getting a patent and attendance at a lecture by Dr. John Mauchly, of ENIAC fame. In 1940, Atanasoff attended a lecture given by Dr. Mauchly and, after the lecture, spent some time discussing electronic digital computers with him. This led to an offer to show Mauchly the ABC; it was later claimed that Mauchly used many of Atanasoff's ideas in the design of the ENIAC without giving proper credit.

Atanasoff sued Mauchly in U.S. Federal Court, charging him with piracy of intellectual property. This trial was not concluded until 1972, at which time U.S. District Judge Earl R. Larson ruled that the ENIAC was "derived" from the ideas of Dr. Atanasoff. While Mauchly was not deemed to have stolen Atanasoff's ideas, the judge did give Atanasoff credit as the inventor of the electronic digital computer. [R32]

As we shall see below, the claim of the ENIAC as the first electronic digital computer is not completely unfair to Dr. Atanasoff; it is just not entirely true.

The Electronic Numerical Integrator and Calculator (ENIAC)

The next significant development in computing machines took place at Moore School of Engineering at the University of Pennsylvania. The ENIAC, designed by J. Presper Eckert and John Mauchly, became operational in 1945.

At the time the ENIAC was proposed, it was well understood that such a machine could be built. The problem was the same as noted by Zuse and his assistants – how to create a large machine using the fragile vacuum tubes of the day. The ENIAC, as finally built, had over 18,000 vacuum tubes. Without special design precautions, such a complex electronic machine would not have remained functional for any significant time interval, let alone perform any useful work.

The steps taken by Eckert to solve the tube reliability problem are described in a paper “Classic Machines: Technology, Implementation, and Economics” in [R11]

“J. Presper Eckert solved the tube reliability problem with a number of techniques:

- Tube failures were bimodal, with more failures near the start of operation as well as after many hours and less during the middle operating hours. In order to achieve better reliability, he used tubes that had already been burnt-in [powered up and left to run for a short time], and had not suffered from ‘infant mortality’ [had not burned out when first turned on].
- The machine was not turned off at night because power cycling [turning the power off and on] was a source of increased failures.
- Components were run at only half of their rated capacity, reducing the stress and degradation of the tubes and significantly increasing their operating lifetime.
- Binary circuits were used, so tubes that were becoming weak or leaky would not cause an immediate failure, as they would in an analog computer.
- The machine was built from circuit modules whose failure was easy to ascertain and that could be quickly replaced with identical modules.”

“As a result of this excellent engineering, ENIAC was able to run for as long as a day or two without a failure! ENIAC was capable of 5,000 operations per second. It took up 2,400 cubic feet [about 70 cubic meters] of circuitry, weighed 30 tons, and consumed 140kW of power.”

While the ENIAC lacked what we would call memory, it did have twenty ten-digit accumulators used to store intermediate results. Each of the ten digits in an accumulator was stored in a shift register of ten flip-flops (for a total of 100 flip-flops per register). According to our reference, a 10–digit accumulator required 550 vacuum tubes, so that the register file by itself required 11,000 vacuum tubes. A bit of arithmetic will show that a 10–digit integer requires 34 bits to represent. Roughly speaking, each accumulator held four 8–bit bytes.

It is this arrangement of the accumulators that might give rise to some confusion. In its instruction architecture, the ENIAC was a decimal machine and not a binary machine; each number being stored as a collection of digits rather than in binary form. As noted above, the

ENIAC did use binary devices to store those digits, probably coded with BCD code, and thus gained reliability.

The next figure shows a module for a single digit in a decimal accumulator. Although this circuit was not taken from the ENIAC, it is sufficiently compatible with the technology used in the ENIAC that it might have been taken from that machine.



Single Digit from a Burroughs 205 Accumulator, circa 1954.

Source: University of Virginia [R33]

The ENIAC was physically a very large machine as shown by the pictures below.

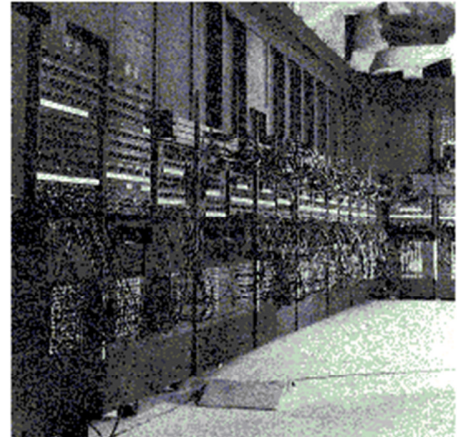


Figure: Two Pictures of the ENIAC

Source: R34

There are several ways to describe the size of the ENIAC. The pictures above, showing an Army private standing at an instrument rack is one. Other pictures include two women programming the ENIAC (using plug cables) or four women, each of whom is holding the circuit to store a single digit in a progression of technologies.

Miss Patsy Simmers,	holding an ENIAC board	(Operational in Fall 1945)
Mrs. Gail TAYlor,	holding an EDVAC board	(Delivered in August 1949)
Mrs. Milly Beck,	holding an ORDVAC board	(Operational in March 1952)
Mrs. Norma Stec,	holding a BRLESC-I board	(Delivered in May 1961)



Figure: Four Women With Register Circuits

It is interesting to note that in 1997, as a part of the 50th anniversary of the ENIAC, a group of students at the University of Pennsylvania created the “ENIAC™-on-a-chip”. It was fabricated in a 0.5-micron, triple metal, CMOS processor. The chip had a size of 7.44 by 5.29 millimeters and contained about 175,000 transistors. [R35]



Figure: Two women programming the ENIAC
(Miss Gloria Ruth Gorden on the left and Mrs. Ester Gertson on the right)

Some Reminiscences about the ENIAC

The following is taken from an interview with J. Presper Eckert, one of the designers of the ENIAC. It was published in the trade journal, Computerworld, on February 20, 2006.

“In the early 1940’s, J. Presper Eckert was the designer and chief engineer building ENIAC, the first general-purpose all-electronic computer (see story, page 18). It was a huge undertaking; ENIAC was the largest electronic device that had ever been built. So why did Eckert – on a tight schedule and with a limited staff – take time out to feed electrical wire to mice?

Because he knew that ENIAC’s hundreds of miles of wiring would be chewed by the rodents. So he used a cageful of mice to taste-test wire samples. The wire whose insulation the mice chewed on least was the stuff Eckert’s team used to wire up the ENIAC. It was an elegant solution to an unavoidable problem.”

In the same issue, an earlier interview with Mr. Eckert was recalled.

“In our 1989 interview, I asked ENIAC co-inventor J. Presper Eckert to recall the zaniest thing he did while developing the historic computer. His reply:

‘The mouse cage was pretty funny. We knew mice would eat the insulation off the wires, so we got samples of all the wires that were available and put them in a cage with a bunch of mice to see which insulation they did not like. We used only wire that passed the mouse test’.”

The Electronic Discrete Variable Automatic Calculator (EDVAC)

The ENIAC, as a result of its quick production in response to war applications, suffered from a number of shortcomings. The most obvious shortcoming is that it had no random access memory. As it had no place to store a program, it was not a stored-program computer.

The EDVAC, first proposed in 1945 and made operational in 1951, was the first true stored program computer in that it had a memory that could be used to store the program. In this sense, the EDVAC was the earliest design that we would call a computer. This design arose from the experience gained on the ENIAC and was proposed in a report “First Draft of a Report on the EDVAC” written in June 1945 by J. Presper Eckert, John Mauchly, and John von Neumann. The report was distributed citing only von Neumann as the author, so that it has only been recently that Eckert and Mauchly have been recognized as coauthors. In the meantime, the name “von Neumann machine” has become synonymous with “stored program computer” as if von Neumann had created the idea all by himself.

Unlike modern computers, the EDVAC and EDSAC (described below) used bit serial arithmetic. This choice was driven by the memory technology available at the time; it was a serial memory based on mercury acoustic delay lines; each bit of a word had to be either read or written sequentially. To access a given word (44 bits) on its memory, the EDVAC had to wait on the given word to circulate to the memory reader and then read it out, one bit at a time. This memory held 1,024 44-bit words, about 5.5 kilobytes.

The EDVAC had almost 6,000 vacuum tubes and 12,000 solid-state diodes. The CPU addition time was 864 microseconds; the multiplication time 2.9 milliseconds. On average, the EDVAC had a mean time between failures of eight hours and ran 20 hours a day.

The goals of the EDVAC, as quoted from the report, were as follows. Note the use of the word “organs” for what we would now call “components” or “circuits”. [R 37]

- “1. Since the device is primarily a computer it will have to perform the elementary operations of arithmetic most frequently. Therefore, it should contain specialized organs for just these operations, i.e., addition, subtraction, multiplication, and division.”
- “2. The logical control of the device (i.e., proper sequencing of its operations) can best be carried out by a central control organ.”
- “3. A device which is to carry out long and complicated sequences of operation must have a considerable memory capacity.”
- “4. The device must have organs to transfer information from the outside recording medium of the device into the central arithmetic part and central control part, and the memory. These organs form its input.”
- “5. The device must have organs to transfer information from the central arithmetic part and central control part, and the memory into the outside recording medium. These organs form its output.”

At this point the main challenge in the development of stored program computers is the availability of technologies suitable for creation of the memory. The vacuum tubes available at the time could have been used to create modest-sized memories, but the memory would have been too expensive and not reliable. Consider an 8KB memory – 8,192 bytes of 8 bits each, corresponding to 65,536 bits. A vacuum tube flip-flop implementation of such a memory would have required one vacuum tube per bit, for a total of 65,536 vacuum tubes. While the terms here are anachronistic, the problem was very real for that time.

Given the above calculation, we see that the big research issue for computer science in the late 1940's and early 1950's was the development of inexpensive and reliable memory elements. Early attempts included electrostatic memories and various types of acoustic delay lines (such as memory delay lines).

The Electronic Delay Storage Automatic Calculator (EDSAC)

While the EDVAC stands as the first design for a stored program computer, the first such computer actually built was the EDSAC by Maurice Wilkes of Cambridge University. The EDSAC became operational on May 6, 1949; the EDVAC in 1951.

In building the EDSAC, the first problem that Wilkes attacked was that of devising a suitable memory unit. By early 1947, he had designed and built the first working mercury acoustic delay memory (see the following figure). It consisted of 16 steel tubes, each capable of storing 32 words of 17 bits each (16 bits plus a sign) for a total of 512 words. The delay lines were bit-serial memories, so the EDSAC had a bit-serial design. The reader is invited to note that the EDSAC was a very large computer, but typical of the size for its time.

The next two figures illustrate the EDSAC I computer, as built at Cambridge.

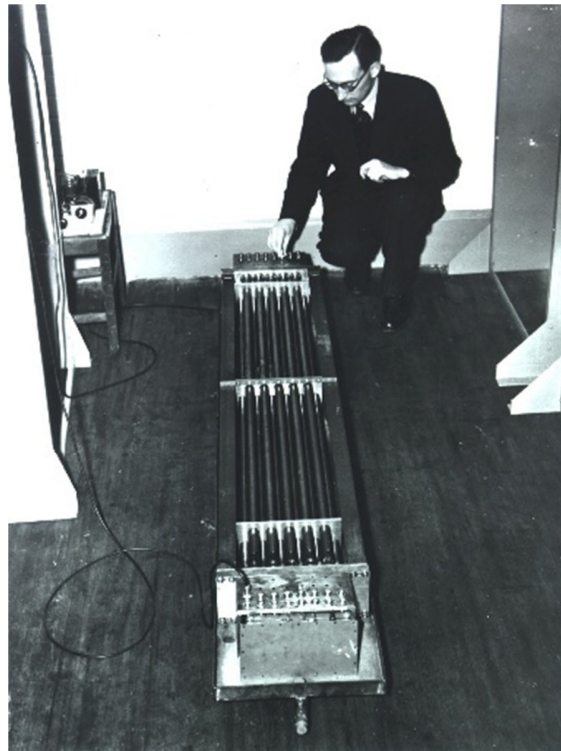


Figure: Maurice Wilkes and a Set of 16 mercury delay lines for the EDSAC [R 38]

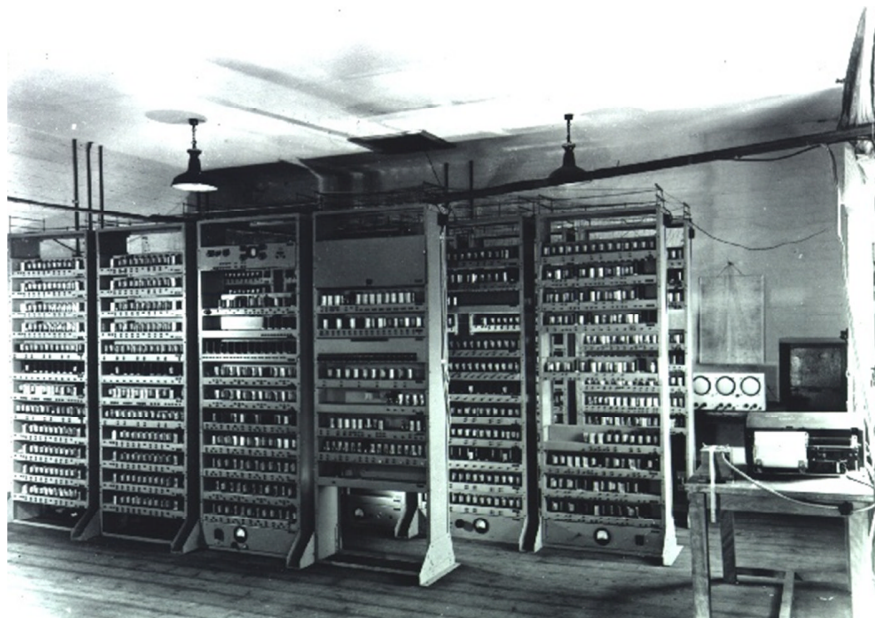


Figure: The EDSAC I [R39]

Note that the computer was built large, with components on racks, so that technicians could have easy access to the electronics when doing routine maintenance.

The design of the EDSAC was somewhat conservative for the time, as Wilkes wanted to create a machine that he could actually program. As a result, the machine cycle time was two microseconds (500 kHz clock rate), slower by a factor of two than other machines being designed at the time. The project pioneered a number of design issues, including the first use of a subroutine call instruction that saved the return address and the use of a linking loader. The subroutine call instruction was named the Wheeler jump, after David J. Wheeler who invented it while working as a graduate student [R11, page 3]. Another of Wheeler's innovations was the relocating program loader.

Magnetic Core Memory

The reader will note that most discussions of the “computer generations” focus on the technology used to implement the CPU (Central Processing Unit). Here, we shall depart from this “CPU-centric view” and discuss early developments in memory technology. We shall begin this section by presenting a number of the early technologies and end it with what your author considers to be the last significant contribution of the first generation – the magnetic core memory. This technology was introduced with the Whirlwind computer, designed and built at MIT in the late 1940's and early 1950's.

The reader will recall that the ENIAC did not have any memory, but just used twenty 10-digit accumulators. For this reason, it was not a stored program computer, as it had no memory in which to store the program. To see the reason for this, let us perform a thought experiment and design a 4 kilobyte (4,096 byte) memory for the ENIAC. We shall imagine the use of flip-flops fabricated from vacuum tubes, a technology that was available to the developers of the ENIAC. We imagine the use of dual-triode tubes, also readily available and well understood during the early and late 1940's.

A dual-triode implementation requires one tube per bit stored, or eight tubes per byte. Our four kilobyte memory would thus require 32,768 tubes just for the data storage. To this, we would add a few hundred vacuum tubes to manage the memory; about 34,000 tubes in all.

Recall that the original ENIAC, having 18 thousand tubes, consumed 140 kilowatts of power might run one or two days without a failure. Our new 52 thousand tube design would have cost an exorbitant amount, consumed 400 kilowatts of power, and possibly run 6 hours at a time before a failure due to a bad tube. An ENIAC with memory was just not feasible.

The EDSAC

Maurice Wilkes of Cambridge University in England was the first to develop a working memory for a digital computer. By early 1947 he had designed and built a working memory based on mercury delay lines (see the picture above). While not large, this memory did allow Wilkes to begin experimentation with a stored program computer and, as indicated above, lead to the development of a number of important programming techniques. However, it was obvious even then that another memory technology was needed if computers were to be developed and used commercially.

The Manchester Mark I

The next step in the development of memory technology was taken at the University of Manchester in England, with the development of the Mark I computer and its predecessor, the “Manchester Baby”. These computers were designed by researchers at the University; the manufacture of the Mark I was contracted out to the Ferranti corporation.

Each of the Manchester Baby and the Manchester/Ferranti Mark I computers used a new technology, called the “Williams–Kilburn tube”, as an electrostatic main memory. This device was developed in 1946 by Frederic C. Williams and Tom Kilburn [R43]. The first design could store either 32 or 64 words (the sources disagree on this) of 32 bits each.

The tube was a cathode ray tube in which the binary bits stored would appear as dots on the screen. One of the early computers with which your author worked was equipped with a Williams–Kilburn tube, but that was no longer in use at the time (1972). The next figure shows the tube, along with a typical display; this has 32 words of 40 bits each.

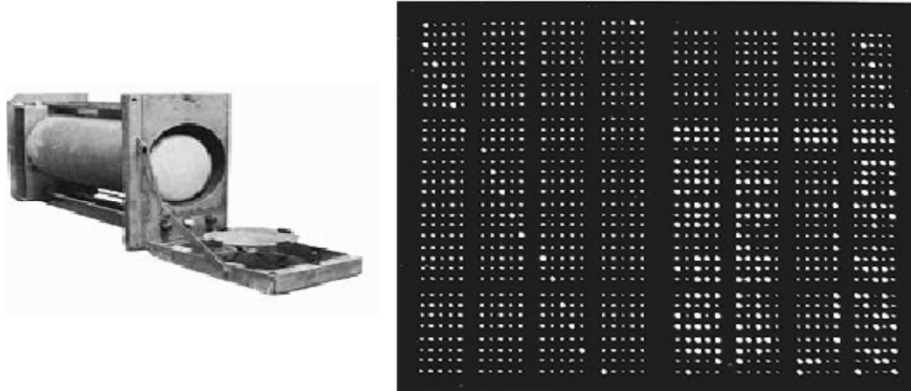


Figure: An Early Williams–Kilburn tube and its display
Note the extra line of 20 bits at the top.

While the Williams–Kilburn tube sufficed for the storage of the prototype Manchester Baby, it was not large enough for the contemplated Mark I. The design team decided to build a hierarchical memory, using double-density Williams–Kilburn tubes (each storing two pages of thirty-two 40-bit words), backed by a magnetic drum holding 16,384 40-bit words.

The magnetic drum can be seen to hold 512 pages of thirty-two 40-bit words each. The data were transferred between the drum memory and the electrostatic display tube one page at a time, with the revolution speed of the drum set to match the display refresh speed. The team used the experience gained with this early block transfer to design and implement the first virtual memory system on the Atlas computer in the late 1950's.

Here is a picture of the drum memory from the IBM 650, from about 1954. This drum was 4 inches in diameter, 16 inches long, and stored 2,000 10-digit numbers. When in use, the drum rotated at 12,500 rpm, or about 208 revolutions per second.



Figure: The IBM 650 Drum Memory

The MIT Whirlwind and Magnetic Core Memory

The MIT Whirlwind was designed as a bit-parallel machine with a 5 MHz clock. The following quote describes the effect of changing the design from one using current technology memory modules with the new core memory modules.

“Initially it [the Whirlwind] was designed using a rank of sixteen Williams tubes, making for a 16-bit parallel memory. However, the Williams tubes were a limitation in terms of operating speed, reliability, and cost.”

“To solve this problem, Jay W. Forrester and others developed magnetic core memories. When the electrostatic memory was replaced with a primitive core memory, the operating speed of the Whirlwind doubled, and the maintenance time on the machine fell from four hours per day to two hours per week, and the mean time between memory failures rose from two hours to two weeks! Core memory technology was crucial in enabling Whirlwind to perform 50,000 operations per second in 1952 and to operate reliably for relative long periods of time” [R11]

The reader should consider the previous paragraph carefully and ask how useful computers would be if we had to spend about half of each working day in maintaining them.

Basically core memory consists of magnetic cores, which are small doughnut-shaped pieces of magnetic material. These can be magnetized in one of two directions: clockwise or counter-clockwise; one orientation corresponding to a logic 0 and the other to a logic 1. Associated with the cores are current carrying conductors used to read and write the memory. The individual cores were strung by hand on the fine insulated wires to form an array. It is a fact that no practical automated core-stringing machine was ever developed; thus core memory was always a bit more costly due to the labor intensity of its construction.

The original core memory modules were quite large. This author used a Digital Equipment Corporation PDP-9 in 1972 while doing his Ph.D. research at Vanderbilt University. The machine had 8K of 18-bit words, placed in two core memory modules. Each module was approximately the size of a 25” by 20” air filter as commonly used in heating systems.

To more clearly illustrate core memory, we include picture from the Sun Microsystems web site. Note the cores of magnetic material with the copper wires interlaced.

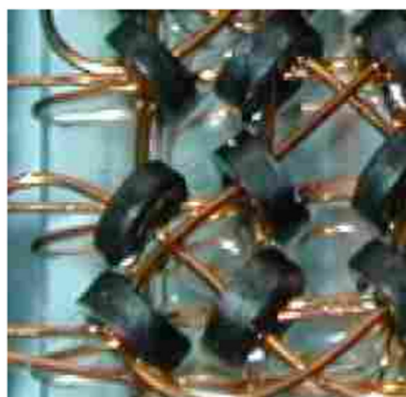


Figure: Close-Up of a Magnetic Core Memory [R40]

The IBM System/360 Model 85 and Cache Memory

Although the development of magnetic core memory represented a significant development in computer technology, it did not solve one of the essential problems. There are two goals for the development of a memory technology; it should be cheap and it should be fast. In general, these two goals are not compatible.

The problem of providing a large inexpensive memory that appeared to be fast had been considered several times before it was directly addressed on the IBM System/360, Model 85. The System/360 family of third-generation computers will be discussed in some detail in the next pages; here we want to focus on its contribution to memory technology. Here we quote from the paper by J. S. Liptay [R58], published in 1968. In that paper, Liptay states:

“Among the objectives of the Model 85 is that of providing a System/360 compatible processor with both high performance and high throughput. One of the important ingredients of high throughput is a large main storage capacity. However, it is not feasible to provide a large main storage with an access time commensurate with the 80-nanosecond processor cycle of the Model 85. ... [W]e decided to use a storage hierarchy. The storage hierarchy consists of a 1.04-microsecond main storage and a small, fast store called a cache, which is integrated into the CPU. The cache is not addressable by a program, but rather is used to hold those portions of main storage that are currently being used.”

The main storage unit for the computer was either an IBM 3265-5 or an IBM 2385, each of which had a capacity between 512 kilobytes and 4 megabytes. The cache was a 16 kilobyte integrated storage, capable of operating at the CPU cycle time of 80 nanoseconds. Optional upgrades made available a cache with either 24 kilobytes or 32 kilobytes.

Both the cache and main memory were divided into sectors, each holding 1 kilobyte of address space that was aligned on 1 KB address boundaries. Each sector was further divided into sixteen blocks of 64 bytes each; these correspond to what we now call “pages”. The cache memory was connected to the main memory via a 16-byte (128 bit) data path so that the processor could fetch or store an entire block in four requests to the main memory. The larger main memory systems were four-way interleaved, so that a request for a block was broken into four requests for 16 bytes each, and sent to each of the four memory banks.

While it required two processor cycles to fetch data in the cache (one to check for presence and one to fetch the data), it was normally the case that requests could be overlapped. In that case, the CPU could access data on every cycle.

In order to evaluate various cache management algorithms, the IBM engineers compared real computers to a hypothetical computer having a single large memory operating at cache speed. This was seen as the upper limit for the real Model 85. For a selected set of test programs, the real machine was seen to operate at an average of 81% of the theoretical peak performance; presumably meaning that it acted as if it had a single large memory with access time equal to 125% that of the cache access time.

Almost all future computers would be designed with a fast cache memory fronting a larger and slow memory. Later developments would show the advantage of having a split cache with an I-cache for instructions and a D-cache for data, as well as multilevel caches. We shall cover these advances in later chapters of the textbook; this is where it began

Interlude: Data Storage Technologies In the 1950's

We have yet to address methods commonly used by early computers for storage of off-line data; that is data that were not actually in the main memory of the computer. One of the main methods was the use of 80 column punch cards, such as seen in the figure below.

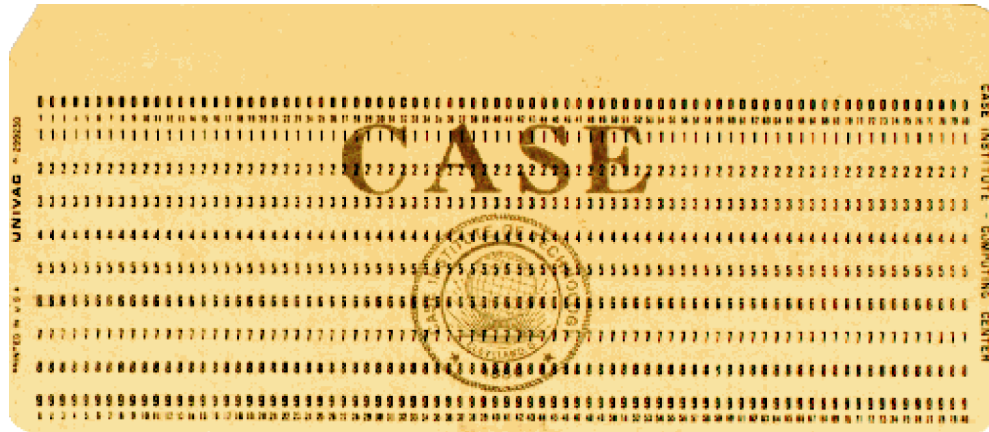


Figure: 80-Column Punch Card (Approximately 75% Size)

These cards would store at most eighty characters, so that the maximum data capacity was exactly eighty bytes. A decent sized program would require a few thousand of these cards to store the source code; these were transported in long boxes. On more than one occasion, this author has witnessed a computer operator drop a box of program cards, scattering the cards in random order and thus literally destroying the computer program.

Seven-Track Tapes and Nine-Track Tapes

Punched cards were widely used by businesses to store account data up to the mid 1950's. At some point, the requirement to keep all of those punch cards became quite burdensome. One life insurance company reported that it had to devote two floors of its modern skyscraper in New York City for the sole purpose of storing data cards. Something had to change.

Data centers began to use larger magnetic tape drives for data storage. Early tape drives were called seven-track units in that they stored seven tracks of information on the tape: six tracks of data and one track of parity bits. Later designs stored eight data bits and parity.



The picture at left shows the first magnetic tape drive produced by IBM. This was the IBM 729, first released in 1952. The novelty of the design is reflected by an actual incident at the IBM receiving dock. The designers were expecting a shipment of magnetic tape when they were called by the foreman of the dock with the news that "We just received a shipment of tape from 3M, but are going to send it back ... It does not have any glue on it"

The next picture shows a typical data center based on the IBM 7090 computer, in use from 1958 through 1969. Note the large number of tape drives. I can count at least 16.



The picky reader will note that the IBM 7090 is a Generation 2 computer that should not be discussed until the next section. We use the picture because the tape drives are IBM 729's, first introduced in 1952, and because the configuration is typical of data centers of the time.

The IBM RAMAC

The next step in evolution of data storage is what became the disk drive. Again, IBM lead the way in developing this storage technology, releasing the RAMAC (Random Access Method for Accounting and Control) in 1956.



Figure: The RAMAC (behind the man's head)

The RAMAC had fifty aluminum disks, each of which was 24 inches in diameter. Its storage capacity is variously quoted at 4.4MB or five million characters. At 5,000,000 characters we have 100,000 characters per 24-inch disk. Not very impressive for data density, but IBM had to start somewhere.

Vacuum Tube Computers (Generation 1 – a parting shot)

By the end of the first generation of computing machines in the middle 1950's, most of the elements of a modern stored program computer had been put in place – the idea of a stored program, a reasonably efficient and reliable memory element, and the basic design of both the central processing unit and the I/O systems.

We leave this generation with a few personal memories of this author. First, there is the IBM 650, a vacuum-tube based computer in production from 1953 through 1962, when the last model was produced. This author recalls joining a research group in 1971, on the day an IBM 650 was being removed to storage to be replaced by a modern PDP-9.



Figure: The IBM 650 – Power Supply, Main Unit, and Read-Punch Unit
Source: Columbia University [R41]

The reader will note that the I/O unit for the computer is labeled a “Read-Punch Unit”. The medium for communicating with this device was the punched card. Thus, the complete system had two components not seen here – a device for punching cards and a device for printing the contents of those cards onto paper for reading by humans.

It was also during the 1950's that this author's Ph.D. advisor, Joseph Hamilton, was a graduate student in a nuclear physics lab that had just received a very new multi-channel analyzer (a variant of a computer). He was told to move it and, unaware that it was top-heavy, managed to tip it over and shatter every vacuum tube in the device. This amounted to a loss of a few thousand dollars – not a good way to begin a graduate career.

One could continue with many stories, including the one about the magnetic drum unit on the first computer put on a sea-going U.S. Navy ship. The drum was spun up and everything was going well until the ship made an abrupt turn and the drum, being basically a large gyroscope, was ripped out of its mount and fell on the floor.

Discrete Transistor Computers (Generation 2 – from 1958 to 1966)

The next generation of computers is based on a device called the transistor. Early forms of transistors had been around for quite some time, but were not regularly used due to a lack of understanding of the atomic structure of crystalline germanium. This author can recall use of an early germanium rectifier in a radio he built as a Cub Scout project – the instructions were to find the “sweet spot” on the germanium crystal for the “cat’s hair” rectifier. Early designs (before the late 1940’s) involved finding this “sweet spot” without any understanding of what made such a spot special. At some time in the 1940’s it was realized that there were two crystalline forms of germanium – called “P type” and “N type”. Methods were soon devised to produce crystals mostly of exactly one of the forms and transistors were possible.

The transistor was invented by Walter Brattain and John Bardeen at Bell Labs in December 1947 and almost instantly improved by William Shockley (also of Bell Labs) in February 1948. The transistor was announced to the world in a press conference during June 1949. In January 1951, Bell labs had worked out most of the difficulties associated with growing the germanium crystals required for transistor manufacture and held a press conference to announce the invention of a working and efficient junction transistor. This was followed in September 1951 by a symposium on transistors attended by about 300 scientists and engineers. Bell announced that it would license the transistor technology to any company willing to use it, and twenty-six companies each paid the fee of \$25,000.

The next year, in April 1952, Bell Labs held the Transistor Technology Symposium. For eight days the attendees worked day and night to learn the tricks of building transistors. At the end of the symposium, they were sent home with a two volume book called “Transistor Technology” that later became known as “Mother Bell’s Cookbook”.

The previous two paragraphs were adapted from material on the Public Broadcast web page <http://www.pbs.org/transistor/>, which contains much additional material on the subject.

While transistors had appeared in early test computers, such as the TX-0 built at MIT, it appears that the first significant computer to use transistor technology was the IBM 7030, informally known as the “Stretch”. Development of the Stretch was begun in 1955 with the goal of producing a computer one hundred times faster than either the IBM 704 or IBM 705, the two vacuum tube computers being sold by IBM at the time, for scientific and commercial use respectively. While missing its goal of being one hundred times faster, the IBM 7030 did achieve significant speed gains by use of new technology. [R11]

1. By switching to germanium transistor technology, the circuitry ran ten times faster.
2. Memory access became about six times faster, due to the use of core memory.

At this time, we pick up our discussion of computer development by noting a series of computers developed and manufactured by the Digital Equipment Corporation of Bedford, Massachusetts. The Digital Equipment Corporation (DEC) had its beginning in the work on test computers at MIT Lincoln Labs in the early 1950’s. In the first years, DEC restricted itself to the manufacture of digital modules, producing its first computer, the PDP-1 (Programmed Data Processor) in 1959.

The technology used in these computers was a module comprised of a small number of discrete transistors and dedicated to one function. The following figure shows a module, called a Flip Chip by DEC, used on the PDP-8 in the late 1960's.

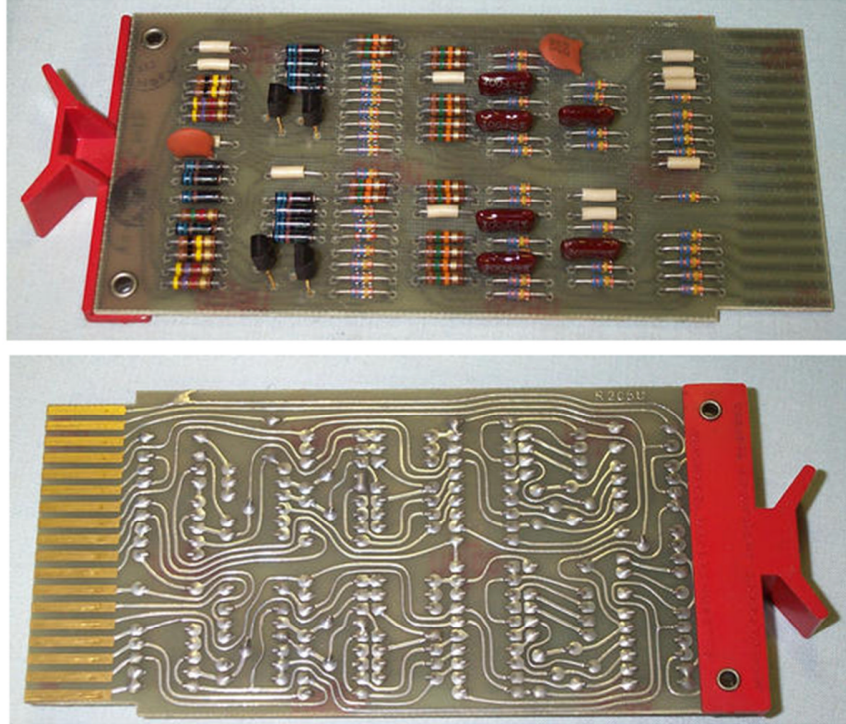


Figure: Front and Back of an R205b Flip-Chip (Dual Flip-Flop) [R42]

In the above figure, one can easily spot many of the discrete components. The orange pancake-like items are capacitors, the cylindrical devices with colorful stripes are resistors with the color coding indicating their rating and tolerances, the black “hats” in the column second from the left are transistors, and the other items must be something useful. The size of this module is indicated by the orange plastic handle, which is fitted to a human hand.

These circuit cards were arranged in racks, as is shown in the next figure.

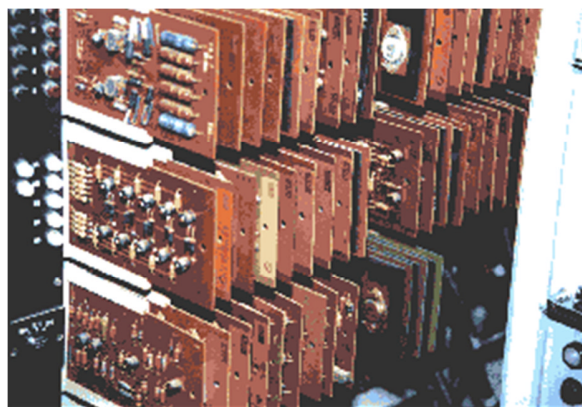


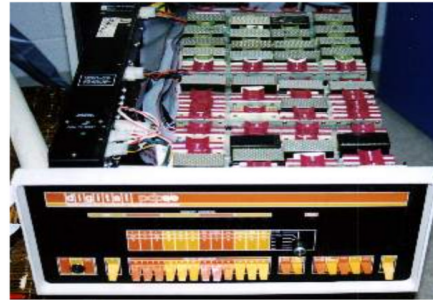
Figure: A Rack of Circuit Cards from the Z-23 Computer (1961)

The PDP-8 was one of the world's first minicomputers, so named because it was small enough to fit on a desk. The name "minicomputer" was suggested by the miniskirt, a knee-length skirt much vogue during that time. This author can say a lot of interesting things about miniskirts of that era, but prefers to avoid controversy.

To show the size of a smaller minicomputer, we use a variant of the PDP-1 from about 1961 and the PDP-8/E (a more compact variant) from 1970. One can see that the PDP-1 with its line of cabinets is a medium-sized computer, while the PDP-8/E is rather small. It is worth note that one of the first video games, Spacewar by Steve Russell, was run on a PDP-1.



PDP-1



PDP-8/E

The history of the development of DEC computers is well documented in a book titled *Computer Engineering: A DEC View of Hardware System Design*, by C. Gordon Bell, J. Craig Mudge, and John E. McNamara [R1]. DEC built a number of "families" of computers, including one line that progressed from the PDP-5 to the PDP-8 and thence to the PDP-12. We shall return to that story after considering the first and second generation computers that lead to the IBM System 360, one of the first third generation computers.

The Evolution of the IBM-360

We now return to a discussion of "Big Iron", a name given informally to the larger IBM mainframe computers of the time. Much of this discussion is taken from an article in the *IBM Journal of Research and Development* [R46], supplemented by articles from [R1]. We trace the evolution of IBM computers from the first scientific computer (the IBM 701, announced in May 1952) through the early stages of the S/360 (announced in March 1964).

We begin this discussion by considering the situation as of January 1, 1954. At the time, IBM has three models announced and shipping. Two of these were the IBM 701 for scientific computations and the IBM 702 for financial calculations (announced in September 1953). Each of the designs used Williams-Kilburn tubes for primary memory, and each was implemented using vacuum tubes in the CPU. Neither computer supported both floating-point (scientific) and packed decimal (financial) arithmetic, as the cost to support both would have been excessive. As a result, there were two "lines": scientific and commercial.

The third model was the IBM 650, mentioned and pictured in the above discussion of the first generation. It was designed as "a much smaller computer that would be suitable for volume production. From the outset, the emphasis was on reliability and moderate cost". The IBM 650 was a serial, decimal, stored-program computer with fixed length words each holding ten decimal digits and a sign. Later models could be equipped with the IBM 305 RAMAC (the disk memory discussed and pictured above). When equipped with terminals, the IBM 650 started the shift towards transaction-oriented processing.

The figure below can be considered as giving the “family tree” of the IBM System/360. Note that there are four “lines”: the IBM 650 line, the IBM 701 line, IBM 702 line, and the IBM 7030 (Stretch). The System/360 (so named because it handled “all 360 degrees of computing”) was an attempt to consolidate these lines and reduce the multiplicity of distinct systems, each with its distinct maintenance problems and costs.

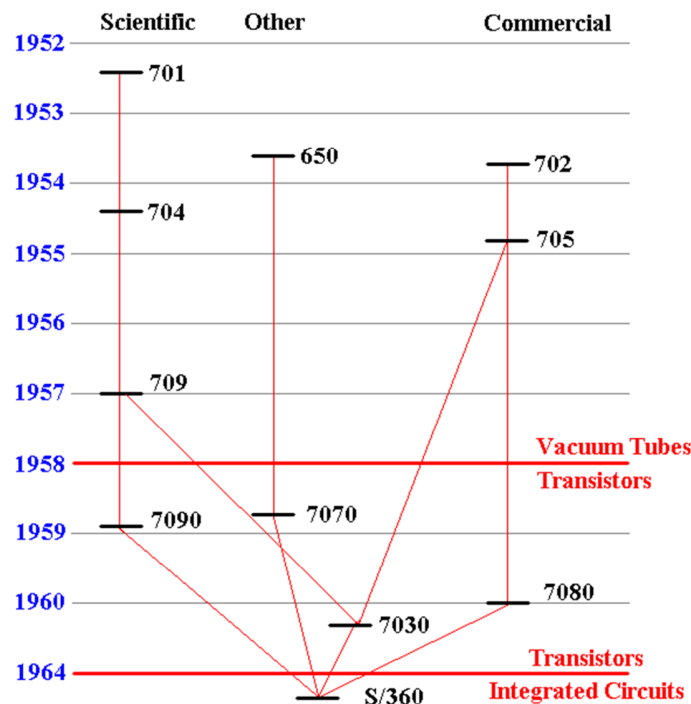


Figure: The IBM System/360 “Family Tree”

As mentioned above, in the 1950’s IBM supported two product lines: scientific computers (beginning with the IBM 701) and commercial computers (beginning with the IBM 702). Each of these lines was redesigned and considerably improved in 1954.

Generation 1 of the Scientific Line

In the IBM 704 (announced in May 1954), the Williams–Kilburn tube memory was replaced by magnetic-core memory with up to 32768 36-bit words. This eliminated the most difficult maintenance problem and allowed users to run larger programs. At the time, theorists had estimated that a large computer would require only a few thousand words of memory. Even at this time, the practical programmers wanted more than could be provided.

The IBM 704 also introduced hardware support for floating-point arithmetic, which was omitted from the IBM 701 in an attempt to keep the design “simple and spartan” [R46]. It also added three index registers, which could be used for loops and address modification. As many scientific programs make heavy use of loops over arrays, this was a welcome addition.

The IBM 709 (announced in January 1957) was basically an upgraded IBM 704. The most important new function was then called a “data-synchronizer unit”; it is now called an “I/O Channel”. Each channel was an individual I/O processor that could address and access main memory to store and retrieve data independently of the CPU. The CPU would interact with the I/O Channels by use of special instructions that later were called channel commands.

It was this flexibility, as much as any other factor, that led to the development of a simple supervisory program called the IOCS (I/O Control System). This attempt to provide support for the task of managing I/O channels and synchronizing their operation with the CPU represents an early step in the evolution of the operating system.

Generation 1 of the Commercial Line

The IBM 702 series differed from the IBM 701 series in many ways, the most important of which was the provision for variable-length digital arithmetic. In contrast to the 36-bit word orientation of the IBM 701 series, this series was oriented towards alphanumeric arithmetic, with each character being encoded as 6 bits with an appended parity check bit. Numbers could have any length from 1 to 511 digits, and were terminated by a “data mark”.

The IBM 705 (announced in October 1954) represented a considerable upgrade to the 702. The most significant change was the provision of magnetic-core memory, removing a considerable nuisance for the maintenance engineers. The size of the memory was at first doubled and then doubled again to 40,000 characters. Later models could be provided with one or more “data-synchronizer units”, allowing buffered I/O independent of the CPU.

Generation 2 of the Product Lines

As noted above, the big change associated with the transition to the second generation is the use of transistors in the place of vacuum tubes. Compared to an equivalent vacuum tube, a transistor offers a number of significant advantages: decreased power usage, decreased cost, smaller size, and significantly increased reliability. These advantages facilitated the design of increasingly complex circuits of the type required by the then new second generation.

The IBM 7070 (announced in September 1958 as an upgrade to the IBM 650) was the first transistor based computer marketed by IBM. This introduced the use of interrupt-driven I/O as well as the SPOOL (Simultaneous Peripheral Operation On Line) process for managing shared Input/Output devices.

The IBM 7090 (announced in December 1958) was a transistorized version of the IBM 709 with some additional facilities. The IBM 7080 (announced in January 1960) likewise was a transistorized version of the IBM 705. Each model was less costly to maintain and more reliable than its tube-based predecessor, to the extent that it was judged to be a “qualitatively different kind of machine” [R46].

The IBM 7090 (and later IBM 7094) were modified by researchers at M.I.T. in order to make possible the CTSS (Compatible Time-Sharing System), the first major time-sharing system. Memory was augmented by a second 32768-word memory bank. User programs occupied one bank while the operating system resided in the other. User memory was divided into 128 memory-protected blocks of 256 words, and access was limited by boundary registers.

The IBM 7094 was announced on January 15, 1962. The CTSS effort was begun in 1961, with a version being demonstrated on an IBM 709 in November 1961. CTSS was formally presented in a paper at the Joint Computer Conference in May, 1962. Its design affected later operating systems, including MULTICS and its derivatives, UNIX and MS-DOS.

As a last comment here, the true IBM geek will note the omission of any discussion of the IBM 1401 line. These machines were often used in conjunction with the 7090 or 7094, handling the printing, punching, and card reading chores for the latter. It is just not possible to cover every significant machine.

The IBM 7030 (Stretch)

In late 1954, IBM decided to undertake a very ambitious research project, with the goal of benefiting from the experience gained in the previous three project lines. In 1955, it was decided that the new machine should be at least 100 times as fast as either the IBM 704 or the IBM 705; hence the informal name “Stretch” as it “stretched the technology”.

In order to achieve these goals, the design of the IBM 7030 required a considerable number of innovations in technology and computer organization; a few are listed here.

1. A new type of germanium transistor, called “drift transistor” was developed. These faster transistors allowed the circuitry in the Stretch to run ten times faster.
2. A new type of core memory was developed; it was 6 times faster than the older core.
3. Memory was organized into multiple 128K-byte units accessed by low-order interleaving, so that successive words were stored in different banks. As a result, new data could be retrieved at a rate of one word every 200 nanoseconds, even though the memory cycle time was 2.1 microseconds (2,100 nanoseconds).
4. Instruction lookahead (now called “pipelining”) was introduced. At any point in time, six instructions were in some phase of execution in the CPU.
5. New disk drives, with multiple read/write arms, were developed. The capacity and transfer rate of these devices lead to the abandonment of magnetic drums.
6. A pair of boundary registers were added so as to provide the storage protection required in a multiprogramming environment.

It is generally admitted that the Stretch did not meet its design goal of a 100 times increase in the performance of the earlier IBM models. Here is the judgment by IBM from 1981 [R46].

“For a typical 704 program, Stretch fell short of its performance target of one hundred times the 704, perhaps by a factor of two. In applications requiring the larger storage capacity and word length of Stretch, the performance factor probably exceeded one hundred, but comparisons are difficult because such problems were not often tackled on the 704.” [R46]

It seems that production of the IBM 7030 (Stretch) was limited to nine machines, one for Los Alamos National Labs, one (called “Harvest”) for the National Security Agency, and 7 more.

Interlude: The Evolution of the PDP-8

While this story does not exactly fit here, it serves as a valuable example of the transition from the second generation of computers to later generations. We quickly trace the CPU size of the PDP-8 computer by the Digital Equipment Corporation from its predecessor (the PDP-5, introduced in 1963), through the first PDP-8 (1964), to a later implementation.

According to Bell [R1], the first PDP-5, built in 1963, had only two registers, an accumulator and the program counter, with the program counter (PC) being stored in memory as word 0. The computer was fabricated from 100 logic boards with approximately 2100 square inches of circuit space for the CPU. Figure XX (above) shows typical mounting for logic boards.

In 1971, the PDP-8/E (a smaller version of the PDP-8) had only three boards with 240 square inches of circuit space for the CPU. In 1976, the Intersil Corporation produced a fully-functioning PDP-8 CPU on a single chip, about 1/4 inch on a side. Thus we have a more powerful CPU on a form factor that is reduced in scale by a factor of 33,600.

Smaller Integrated (SSI and MSI) Circuits (Generation 3 – from 1966 to 1972)

In one sense, the evolution of computer components can be said to have stopped in 1958 with the introduction of the transistor; all future developments merely represent the refinement of the transistor design. This statement stretches the truth so far that it hardly even makes this author's point, which is that packaging technology is extremely important.

What we see in the generations following the introduction of the transistor is an aggressive miniaturization of both the transistors and the traces (wires) used to connect the circuit elements. This allowed the creation of circuit modules with component densities that could hardly have been imagined a decade earlier. Such circuit modules used less power and were much faster than those of the second generation; in electronics smaller is faster. They also lent themselves to automated manufacture, thus increasing component reliability.

It seems that pictures are the best way to illustrate the evolution of the first three generations of computer components. Below, we see a picture of an IBM engineer (they all wore coats and ties at that time) with three generations of components.

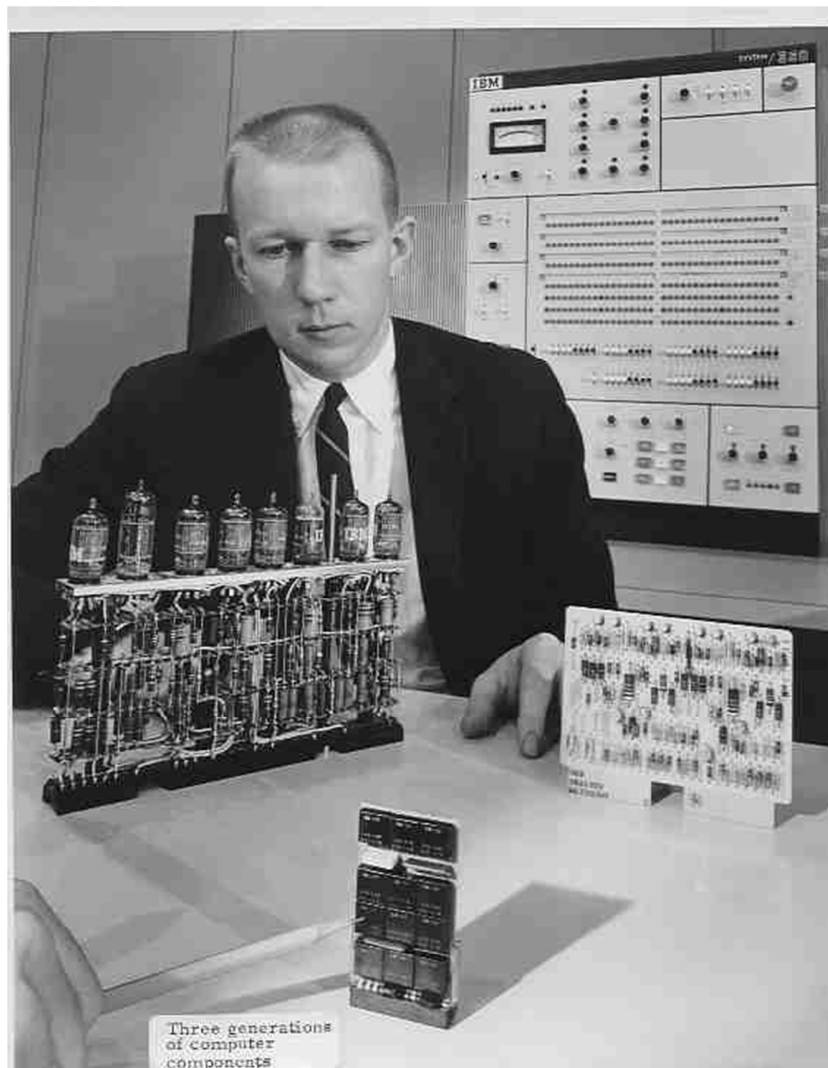


Figure: IBM Engineer with Three Generations of Components

The first generation unit (vacuum tubes) is a pluggable module from the IBM 650. Recall that the idea of pluggable modules dates to the ENIAC; the design facilitates maintenance.

The second generation unit (discrete transistors) is a module from the IBM 7090.

The third generation unit is the ACPX module used on the IBM 360/91 (1964). Each chip was created by stacking layers of silicon on a ceramic substrate; it accommodated over twenty transistors. The chips could be packaged together onto a circuit board. The circuit board in the foreground appears to hold twelve chips. This author conjectures, but cannot prove that the term “ACPX module” refers to the individual chip, and that the pencil in the foreground is indicating one such module.

It is likely that each of the chips on the third-generation board has a processing power equivalent to either of the units from the previous generations.

The 74181 Arithmetic Logic Unit by Texas Instruments

The first real step in the third generation of computing was taken when Texas Instruments introduced the 7400 series of integrated circuits. One of the earliest, and most famous, was a MSI chip called the 74181. It was an arithmetic logic unit (ALU) that provided thirty-two functions of two 4-bit variables, though many of those functions were quite strange. It was developed in the middle 1960's and first used in the Data General Nova computer in 1968.

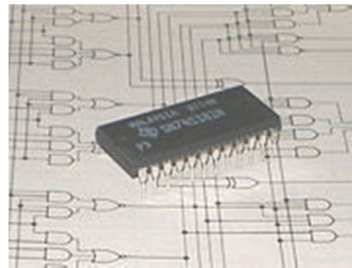


Figure: The 74181 in a DIP Packaging

The figure above shows the chip in its DIP (Dual In-line Pin) packaging. The figure below shows the typical dimensions of a chip in the DIP configuration.

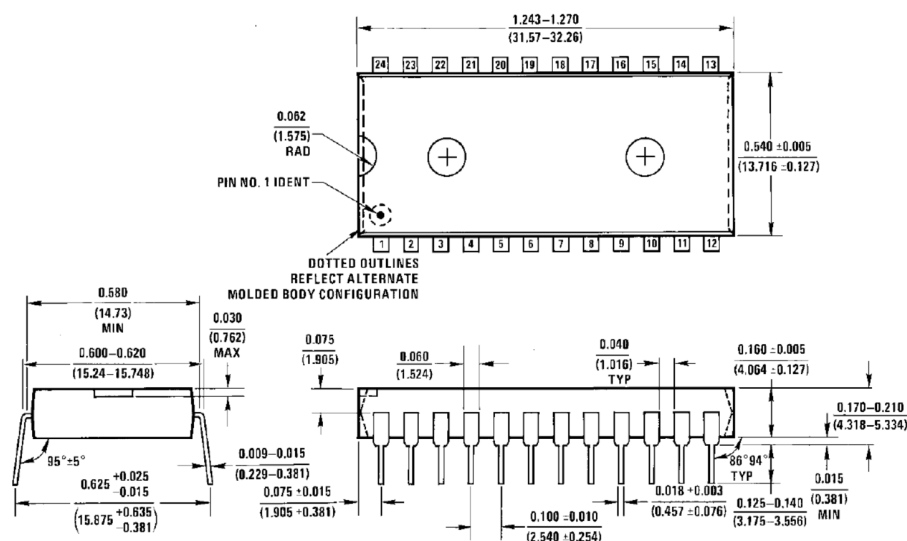
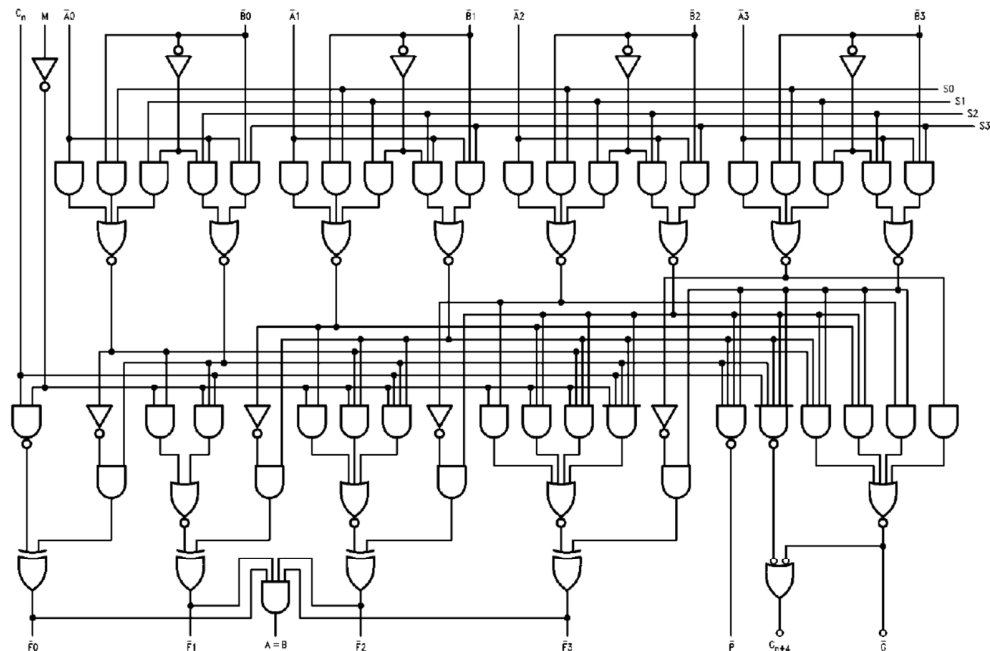


Figure: The 74181 Physical Dimensions in Inches (Millimeters)

For those who like circuit diagrams, here is a circuit diagram of the 74181 chip. A bit later in the textbook, we shall discuss how to read this diagram; for now note that it is moderately complex, comprising seventy five logic gates.

Logic Diagram



Here is a long description of the 74181, taken from the Wikipedia article.

“The 74181 is a 7400 series medium-scale integration (MSI) TTL integrated circuit, containing the equivalent of 75 logic gates and most commonly packaged as a 24-pin DIP. The 4-bit wide ALU can perform all the traditional add / subtract / decrement operations with or without carry, as well as AND / NAND, OR / NOR, XOR, and shift. Many variations of these basic functions are available, for a total of 16 arithmetic and 16 logical operations on two four bit words. Multiply and divide functions are not provided but can be performed in multiple steps using the shift and add or subtract functions.

The 74181 performs these operations on two four bit operands generating a four bit result with carry in 22 nanoseconds. The 74S181 performs the same operations in 11 nanoseconds.

Multiple 'slices' can be combined for arbitrarily large word sizes. For example, sixteen 74S181s and five 74S182 look ahead carry generators can be combined to perform the same operations on 64-bit operands in 28 nanoseconds. Although overshadowed by the performance of today's multi-gigahertz 64-bit microprocessors, this was quite impressive when compared to the sub megahertz clock speeds of the early four and eight bit microprocessors

Although the 74181 is only an ALU and not a complete microprocessor it greatly simplified the development and manufacture of computers and other devices that required high speed computation during the late 1960s through the early 1980s, and is still referenced as a "classic" ALU design.

Prior to the introduction of the 74181, computer CPUs occupied multiple circuit boards and even very simple computers could fill multiple cabinets. The 74181 allowed an entire CPU and in some cases, an entire computer to be constructed on a single large printed circuit board. The 74181 occupies a historically significant stage between older CPUs based on discrete logic functions spread over multiple circuit boards and modern microprocessors that incorporate all CPU functions in a single component. The 74181 was used in various minicomputers and other devices beginning in the late 1960s, but as microprocessors became more powerful the practice of building a CPU from discrete components fell out of favor and the 74181 was not used in any new designs.

Many computer CPUs and subsystems were based on the 74181, including several historically significant models.

1. NOVA - First widely available 16-bit minicomputer manufactured by Data General. This was the first design (in 1968) to use the 74181.
2. PDP-11 - Most popular minicomputer of all time, manufactured by Digital Equipment Corporation. The first model was introduced in 1970.
3. Xerox Alto - The first computer to use the desktop metaphor and graphical user interface (GUI).
4. VAX 11/780 - The first VAX, the most popular 32-bit computer of the 1980s[8], also manufactured by Digital Equipment Corp.”

Back to the System/360

Although the System/360 likely did not use any 74181's, that chip does illustrate the complexity of the custom-fabricated chips used in IBM's design.

The design goals for the System/360 family are illustrated by the following scenario. Imagine that you have a small company that is using an older IBM 1401 for its financial work. You want to expand and possibly do some scientific calculations. Obviously, IBM is very happy to lease you a computer. We take your company through its growth.

1. At first, your company needs only a small computer to handle its computing needs. You lease an IBM System 360/30. You use it in emulation mode to run your IBM 1401 programs unmodified.
2. Your company grows. You need a bigger computer. Lease a 360/50.
3. You hit the “big time”. Turn in the 360/50 and lease a 360/91. You never need to rewrite or recompile your existing programs. You can still run your IBM 1401 programs without modification.

In order to understand how the System/360 was able to address these concerns, we must divert a bit and give a brief description of the control unit of a stored program computer.

The Control Unit and Emulation

In order to better explain one of the distinctive features of the IBM System/360 family, it is necessary to take a detour and discuss the function of the control unit in a stored program computer. Basically, the control unit tells the computer what to do.

All modern stored program computers execute programs that are a sequence of binary machine-language instructions. This sequence of instructions corresponds either to an assembly language program or a program in a higher-level language, such as C++ or Java.

The basic operation of a stored program computer is called “**fetch-execute**”, with many variants on this name. Each instruction to be executed is fetched from memory and deposited in the **Instruction Register**, which is a part of the control unit. The control unit interprets this machine instruction and issues the signals that cause the computer to execute it.

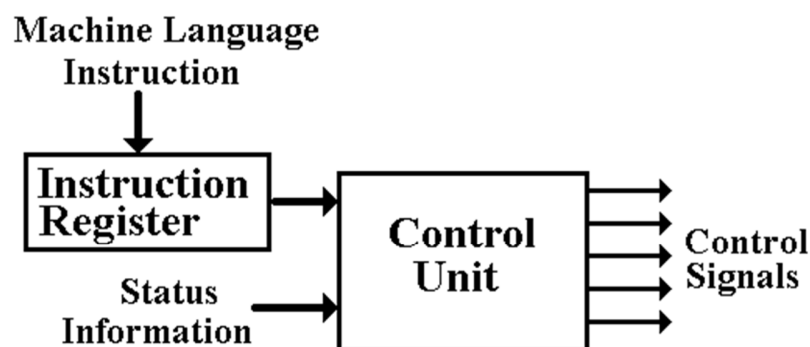


Figure: Schematic of the Control Unit

There are two ways in which a control unit may be organized. The most efficient way is to build the unit entirely from basic logic gates. For a moderately-sized instruction set with the standard features expected, this leads to a very complex circuit that is difficult to test.

In 1951, Maurice V. Wilkes (designer of the EDSAC, see above) suggested an organization for the control unit that was simpler, more flexible, and much easier to test and validate. This was called a “**microprogrammed control unit**”. The basic idea was that control signals can be generated by reading words from a **micromemory** and placing each in an output buffer.

In this design, the control unit interprets the machine language instruction and branches to a section of the micromemory that contains the **microcode** needed to emit the proper control signals. The entire contents of the micromemory, representing the sequence of control signals for all of the machine language instructions is called the **microprogram**. All we need to know is that the microprogram is stored in a ROM (Read Only Memory) unit.

While microprogramming was sporadically investigated in the 1950's, it was not until about 1960 that memory technology had matured sufficiently to allow commercial fabrication of a micromemory with sufficient speed and reliability to be competitive. When IBM selected the technology for the control units of some of the System/360 line, its primary goal was the creation of a unit that was easily tested. Then they got a bonus; they realized that adding the appropriate blocks of microcode could make a S/360 computer execute machine code for either the IBM 1401 or IBM 7094 with no modification. This greatly facilitated upgrading from those machines and significantly contributed to the popularity of the S/360 family.

The IBM System/360

As noted above, the IBM chose to replace a number of very successful, but incompatible, computer lines with a single computer family, the System/360. The goal, according to an IBM history web site [R48] was to “provide an expandable system that would serve every data processing need”. The initial announcement on April 7, 1964 included Models 30, 40, 50, 60, 62, and 70 [R49]. The first three began shipping in mid-1965, and the last three were replaced by the Model 65 (shipped in November 1965) and Model 75 (January 1966).

The introduction of the System/360 is also the introduction of the term “architecture” as applied to computers. The following quotes are taken from one of the first papers describing the System/360 architecture [R46].

“The term *architecture* is used here to describe the attributes of a system as seen by the programmer,, i.e., the conceptual structure and functional behavior, as distinct from the organization of the data flow and controls, the logical design, and the physical implementation.”

“In the last few years, many computer architects have realized, usually implicitly, that logical structure (as seen by the programmer) and physical structure (as seen by the engineer) are quite different. Thus, each may see registers, counters, etc., that to the other are not at all real entities. This was not so in the computers of the 1950’s. The explicit recognition of the duality of the structure opened the way for the compatibility within System/360.”

At this point, we differentiate between the three terms: architecture, organization, and implementation. The quote above gives a precise comparison of the terms architecture and organization. As an example, consider the sixteen general-purpose registers (R0 – R15) in the System/360 architecture. All models implement these registers and use them in exactly the same way; this is a requirement of the **architecture**. As a matter of **implementation**, a few of the lower end S/360 family actually used dedicated core memory for the registers, while the higher end models used solid state circuitry on the CPU board.

The difference between organization and implementation is seen by considering the two computer pairs: the 709 and 7090, and the 705 and 7080. The IBM 7090 had the same organization (hence the same architecture) as the IBM 709; the implementation was different. The IBM 709 used vacuum tubes; the IBM 7090 replaced these with transistors.

The requirement for the System/360 design is that all models in that series would be “strictly program compatible, upward and downward, at the program bit level”. [R46]

“Here it [strictly program compatible] means that a valid program, whose logic will not depend implicitly upon time of execution and which runs upon configuration A, will also run on configuration B if the latter includes as least the required storage, at least the required I/O devices, and at least the required optional features.”

“Compatibility would ensure that the user’s expanding needs be easily accommodated by any model. Compatibility would also ensure maximum utility of programming support prepared by the manufacturer, maximum sharing of programs generated by the user, ability to use small systems to back up large ones, and exceptional freedom in configuring systems for particular applications.”

Additional design goals for the System/360 include the following.

1. The System/360 was intended to replace two mutually incompatible product lines in existence at the time.
 - a) The scientific series (701, 704, 7090, and 7094) that supported floating point arithmetic, but not decimal arithmetic.
 - b) The commercial series (702, 705, and 7080) that supported decimal arithmetic, but not floating point arithmetic.
2. The System/360 should have a “compatibility mode” that would allow it to run unmodified machine code from the IBM 1401 – at the time a very popular business machine with a large installed base.

This was possible due to the use of a microprogrammed control unit. If you want to run native S/360 code, access that part of the microprogram. If you want to run IBM 1401 code, just switch to the microprogram for that machine.
3. The Input/Output Control Program should be designed to allow execution by the CPU itself (on smaller machines) or execution by separate I/O Channels on the larger machines.
4. The system must allow for autonomous operation with very little intervention by a human operator. Ideally this would be limited to mounting and dismounting magnetic tapes, feeding punch cards into the reader, and delivering output.
5. The system must support some sort of extended precision floating point arithmetic, with more precision than the 36-bit system then in use.

Minicomputers: the Rise and Fall of the VAX

So far, our discussion of the third generation of computing has focused on the mainframe, specifically the IBM System/360. Later, we shall mention another significant large computer of this generation; the CDC-6600. For now, we turn our attention to the minicomputer, which represents another significant development of the third generation. Simply, these computers were smaller and less costly than the “big iron”.

In 1977, Gordon Bell provided a definition of a minicomputer [R1, page 14].

“[A minicomputer is a] computer originating in the early 1960s and predicated on being the lowest (minimum) priced computer built with current technology. From this origin, at prices ranging from 50 to 100 thousand dollars, the computer has evolved both at a price reduction rate of 20 percent per year ...”

The above definition seems to ignore microcomputers; probably these were not “big players” at the time it was made. From the modern perspective, the defining characteristics of a minicomputer were relatively small size and modest cost.

As noted above, the PDP-8 by Digital Equipment Corporation is considered to be the first minicomputer. The PDP-11 followed in that tradition.

This discussion focuses on the VAX series of computers, developed and marketed by the Digital Equipment Corporation (now out of business), and its predecessor the PDP-11. We begin with a discussion of the PDP-11 family, which was considered to be the most popular minicomputer; its original design was quite simple and elegant. We then move to a discussion of the Virtual Address Extension (VAX) of the PDP-11, note its successes, and trace the reasons for its demise. As we shall see, the major cause was just one bad design decision, one that seemed to be very reasonable when it was first made.

The **Programmable Data Processor 11**, introduced in 1970, was an outgrowth of earlier lines of small computers developed by the Digital Equipment Corporation, most notably the PDP-8 and PDP-9. The engineers made a survey of these earlier designs and made a conscious decision to fix what were viewed as design weaknesses. As with IBM and its System/360, DEC decided on a family of compatible computers with a wide range of price (\$500 to \$250,000) and performance. All members of the PDP-11 family were fabricated from TTL integrated circuits; thus the line was definitely third-generation.

The first model in the line was the PDP-11/20, followed soon by the PDP-11/45, a model designed for high performance. By 1978, there were twelve distinct models in the PDP-11 line (LSI-11, PDP-11/04, PDP-11/05, PDP-11/20, PDP-11/34, PDP-11/34C, PDP-11/40, PDP-11/45, PDP-11/55, PDP-11/60, PDP-11/70, and the VAX-11/780). Over 50,000 of the line had been sold, which one of the developers (C. Gordon Bell), considered a qualified success, “since a large and aggressive marketing organization, armed with software to correct architectural inconsistencies and omissions, can save almost any design”[R1, page 379].

With the exception of the VAX-11/780 all members of this family were essentially 16-bit machines with 16-bit address spaces. In fact, the first of the line (the PDP-11/20) could address only $2^{16} = 65536$ bytes of memory. Given that the top 4K of the address space was mapped for I/O devices, the maximum amount of physical memory on this machine was 61,440 bytes. Typically, with the Operating System installed, this left about 32 KB space for user programs. While many programs could run in this space, quite a few could not.

According to two of the PDP-11’s developers:

“The first weakness of minicomputers was their limited addressing capability. The PDP-11 followed this hallowed tradition of skimping on address bits, but it was saved by the principle that a good design can evolve through at least one major change. ... It is extremely embarrassing that the PDP-11 had to be redesigned with memory management [used to convert 16-bit addresses into 18-bit and then 22-bit addresses] only two years after writing the paper that outlined the goal of providing increased address space.”[R1, page 381]

While the original PDP-11 design was patched to provide first an 18-bit address space (the PDP-11/45 in 1972) and a 22-bit address space (the PDP-11/70 in 1975), it was quickly obvious that a new design was required to solve the addressability problem. In 1974, DEC started work on extending the virtual address space of the PDP-11, with the goal of a computer that was compatible with the existing PDP-11 line. In April 1975, work was redirected to produce a machine that would be “culturally compatible with the PDP-11”. The result was the VAX series, the first model of which was the VAX-11/780, which was introduced on October 25, 1977. At the time, it was called a “super-mini”; its later models, such as the VAX-11/9000 series were considered to be mainframes.

By 1982 Digital Equipment Corporation was the number two computer company. IBM maintained its lead as number one. The VAX series of computers, along with its proprietary operating system, VMS, were the company sales leaders. One often heard of “VAX/VMS”.

The VAX went through many different implementations. The original VAX was implemented in TTL and filled more than one rack for a single CPU. CPU implementations that consisted of multiple ECL gate array or macrocell array chips included the 8600, 8800 super-minis and finally the 9000 mainframe class machines. CPU implementations that consisted of multiple MOSFET custom chips included the 8100 and 8200 class machines. The computers in the VAX line were excellent machines, supported by well-designed and stable operating systems. What happened? Why did they disappear?



Figure: A Complete VAX-11/780 System

In order to understand the reason for the demise of the VAX line of computers, we must first present some of its design criteria, other than the expanded virtual address space. Here we quote from Computer Architecture: A Quantitative Approach by Hennessy & Patterson.

“In the late 1960’s and early 1970’s, people realized that software costs were growing faster than hardware costs. In 1967, McKeeman argued that compilers and operating systems were getting too big and too complex and taking too long to develop. Because of inferior compilers and the memory limitations of computers, most system programs at the time were still written in assembly language. Many researchers proposed alleviating the software crisis by creating more powerful, software-oriented architectures.

In 1978, Strecker discussed how he and the other architects at DEC responded to this by designing the VAX architecture. The VAX was designed to simplify compilation of high-level languages. ... The VAX architecture was designed to be highly orthogonal and to allow the mapping of high-level-language statements into single VAX instructions.” [R60, page 126]

What we see here is the demand for what would later be called a CISC design. In a Complex Instruction Set Computer, the ISA (Instruction Set Architecture) provides features to support a high-level language, and minimize what was called the “semantic gap”.

The VAX has been perceived as the quintessential CISC processing architecture, with its very large number of addressing modes and machine instructions, including instructions for such complex operations as queue insertion/deletion and polynomial evaluation. Eventually, it was this complexity that caused its demise. As we shall see below, this complexity was not required or even found useful.

The eventual downfall of the VAX line was due to the CPU design philosophy called “RISC”, for Reduced Instruction Set Computer. As we shall see later in the text, there were many reasons for the RISC movement, among them the fact that the semantic gap was more a fiction of the designers imagination than a reality. What compiler writers wanted was a clean CPU design with a large number of registers and possibly a good amount of cache memory. None of these features were consistent with a CISC design, such as the VAX.

The bottom line was that the new RISC designs offered more performance for less cost than the VAX line of computers. Most of the new designs, including those based on the newer Intel CPU chips, were also smaller and easier to manage. The minicomputers that were so popular in the 1970’s and 1980’s had lost their charm to newer and smaller models.

DEC came late to the RISC design arena, marketing the Alpha in 1992, a 64 bit machine. By then it was too late. Most data centers had switched to computers based on the Intel chips, such as the 80486/80487 or the Pentium. The epitaph for DEC is taken from Wikipedia.

“Although many of these products were well designed, most of them were DEC-only or DEC-centric, and customers frequently ignored them and used third-party products instead. Hundreds of millions of dollars were spent on these projects, at the same time that workstations based on RISC architecture were starting to approach the VAX in performance. Constrained by the huge success of their VAX/VMS products, which followed the proprietary model, the company was very late to respond to commodity hardware in the form of Intel-based personal computers and standards-based software such as Unix as well as Internet protocols such as TCP/IP. In the early 1990s, DEC found its sales faltering and its first layoffs followed. The company that created the minicomputer, a dominant networking technology, and arguably the first computers for personal use, did not effectively respond to the significant restructuring of the computer industry.”

“Beginning in 1992, many of DEC’s assets were spun off in order to raise revenue. This did not stem the tide of red ink or the continued lay offs of personnel. Eventually, on January 26, 1998, what remained of the company was sold to Compaq Computer Corporation. In August 2000, Compaq announced that the remaining VAX models would be discontinued by the end of the year. By 2005 all manufacturing of VAX computers had ceased, but old systems remain in widespread use. Compaq was sold to Hewlett-Packard in 2002.”

It was the end of an era; minicomputers had left the scene.

Large Scale and Very Large Scale Integrated Circuits (from 1972 onward)

We now move to a discussion of LSI and VLSI circuitry. We could trace the development of the third generation System/360 through later, more sophisticated, implementations. Rather than doing this, we shall trace the development of another iconic processor series, the Intel 4004, 8008, and those that followed.

The most interesting way to describe the beginning of the fourth generation of technology, that of a microprocessor on a chip [R1], is to quote from preface to the 1995 history of the microcomputer written by Stanley Mazor [R61], who worked for Intel from 1969 to 1984.

“Intel’s founder, Robert Noyce, chartered Ted Hoff’s Applications Research Department in 1969 to find new applications for silicon technology – the microcomputer was the result. Hoff thought it would be neat to use MOS LSI technology to produce a computer. Because of the ever growing density of large scale integrated (LSI) circuits, a ‘computer on a chip’ was inevitable. But in 1970 we could only get about 2000 transistors on a chip and a conventional CPU would need about 10 times that number. We developed two ‘microcomputers’ 10 years ahead of ‘schedule’ by scaling down the requirements and using a few other ‘tricks’ described in this paper.”

Intel delivered two Micro Computer Systems in the early 1970’s. The first was the MCS-4, emphasizing low cost, in November 1971. This would lead to a line of relatively powerful but inexpensive controllers, such as the Intel 8051 (which in quantity sells for less than \$1). The other was the MCS-8, emphasizing versatility, in April 1972. This lead to the Intel 8008, 8080, 8086, 80286, and a long line of processors for personal computers.

The MCS-4 was originally developed in response to a demand for a chip to operate a hand-held calculator. It was a four-bit computer, with four-bit data memory, in response to the use of 4-bit BCD codes to represent the digits in the calculator. The components of the MCS-4 included the 4001 (Read Only Memory), 4002 (Random Access Memory), and the 4004 Microprocessor. Each was mounted in a 16-pin package, as that was the only package format available in the company at the time.

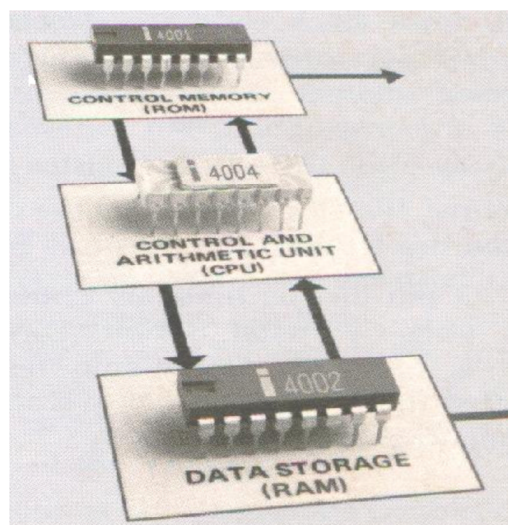


Figure: Picture from the 1972 Spec Sheet on the MCS-4

In 1972, the Intel 4004 sold in quantity for less than \$100 each. It could add two 10-digit numbers in about 800 microseconds. This was comparable to the speed of the IBM 1620, a computer that sold for \$100,000 in 1960.

The MCS-8 was based on the Intel 8008, Intel's first 8-bit CPU. It was implemented with TTL technology, had 48 instructions, and had a performance rating of 0.06 MIPS (Million Instructions per Seconds, a term not in use at the time). It had an 8-bit accumulator and six 8-bit general purpose registers (B, C, D, E, H, and L). In later incarnations of the model, these would become the 8-bit lower halves of the 16-bit registers with the same name.

The Intel 8008 was placed in an 18-pin package. It is noteworthy that the small pin counts available for the packaging drove a number of design decisions, such as multiplexing the bus. In this, the bus connecting the CPU to the memory chip has at least two uses, and a control signal to indicate what signals are on the bus.

The 8008 was followed by the 8080, which had ten times the performance and was housed in a 40-pin plastic package, allowing separate lines for bus address and bus data. The 8080 did not yet directly support 16-bit processing, but arranged its 8-bit registers in pairs. The 8080 was followed by the 8085 and then by the 8086, a true 16-bit CPU. The 8088 (an 8086 variant with a 8-bit data bus) was selected by IBM for its word processor; the rest is history.

As we know, the 8086 was the first in a sequence of processors with increasing performance. The milestones in this line were the 80286, the 80386 (the first with true 32-bit addressing), the 80486/80487, and the Pentium. The line will soon implement 64-bit addressing.

The major factor driving the development of smaller and more powerful computers was and continues to be the method for manufacture of the chips. This is called "photolithography", which uses light to transmit a pattern from a photomask to a light-sensitive chemical called "photoresist" that is covering the surface of the chip. This process, along with chemical treatment of the chip surface, results in the deposition of the desired circuits on the chip.

Due to a fundamental principle of Physics, the minimum feature size that can be created on the chip is approximately the same as the wavelength of the light used in the lithography. State-of-the-art (circa 2007) photolithography uses deep ultraviolet light with wavelengths of 248 and 193 nanometers, resulting in feature sizes as small as 50 nanometers [R53]. An examination of the table on the next page shows that the minimum feature size seems to drop in discrete steps in the years after 1971; each step represents a new lithography process.

One artifact of this new photolithography process is the enormous cost of the plant used to fabricate the chips. In the terminology of semiconductor engineering, such a plant is called a "fab". In 1998, Gordon Moore of Intel Corporation (and Moore's Law fame) noted:

"An R&D fab today costs \$400 million just for the building. Then you put about \$1 billion of equipment in it. That gives you a quarter-micron [250 nanometer] fab for about 5,000 wafers per week, about the smallest practical fab."

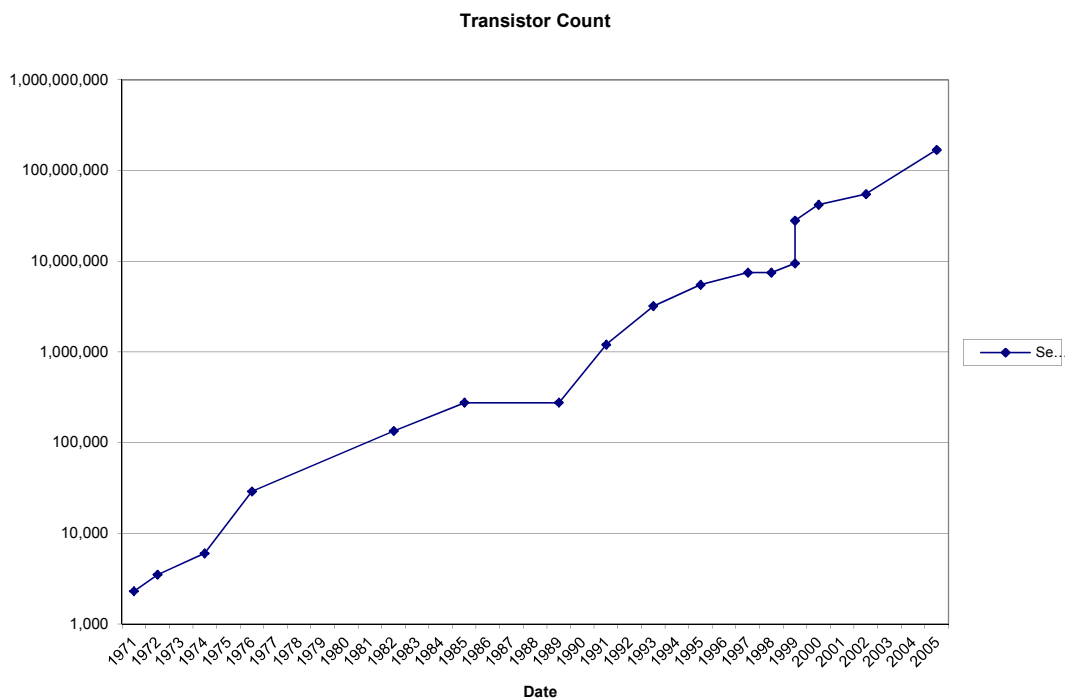
The "quarter-micron fab" represents the ability to create chips in which the smallest line feature has a size of 0.25 micron, or 250 nanometers.

One can describe the VLSI generation of computers in many ways, all of them interesting. However, the most telling description is to give the statistics of a line of computers and let the numbers speak for themselves. The next two pages do just that for the Intel line.

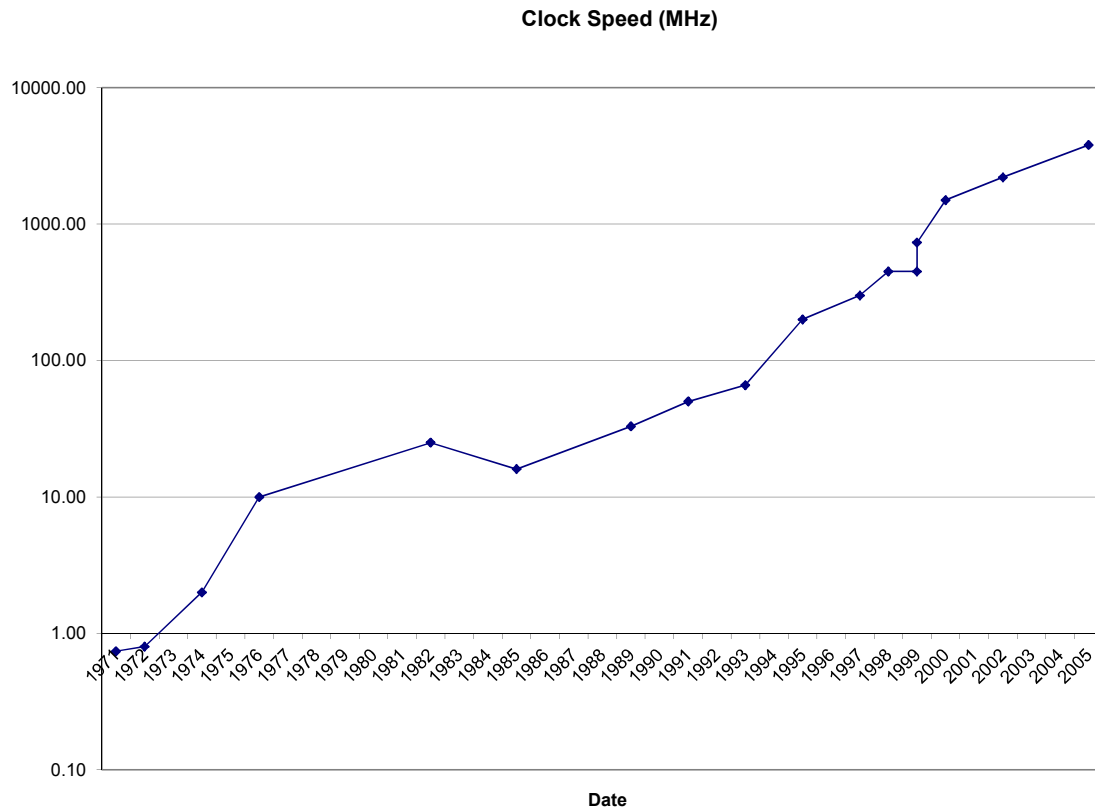
Here are the raw statistics, taken from a Wikipedia article [R52].

Model	Introduced	MHz	Transistors	Line Size
4004	11/15/1971	0.74	2,300	10.00
8008	4/1/1972	0.80	3,500	10.00
8080	4/1/1974	2.00	6,000	6.00
8086	6/8/1976	10.00	29,000	3.00
80286	2/1/1982	25.00	134,000	1.50
80386	10/17/1985	16.00	275,000	1.00
80386	4/10/1989	33.00	275,000	1.00
80486	6/24/1991	50.00	1,200,000	0.80
Pentium	3/23/1993	66.00	3,200,000	0.60
Pentium Pro	11/1/1995	200.00	5,500,000	0.35
Pentium II	5/7/1997	300.00	7,500,000	0.35
Pentium II	8/24/1998	450.00	7,500,000	0.25
Pentium III	2/26/1999	450.00	9,500,000	0.25
Pentium III	10/25/1999	733.00	28,100,000	0.18
Pentium 4	11/20/2000	1500.00	42,000,000	0.18
Pentium 4	1/7/2002	2200.00	55,000,000	0.13
Pentium 4F	2/20/2005	3800.00	169,000,000	0.09

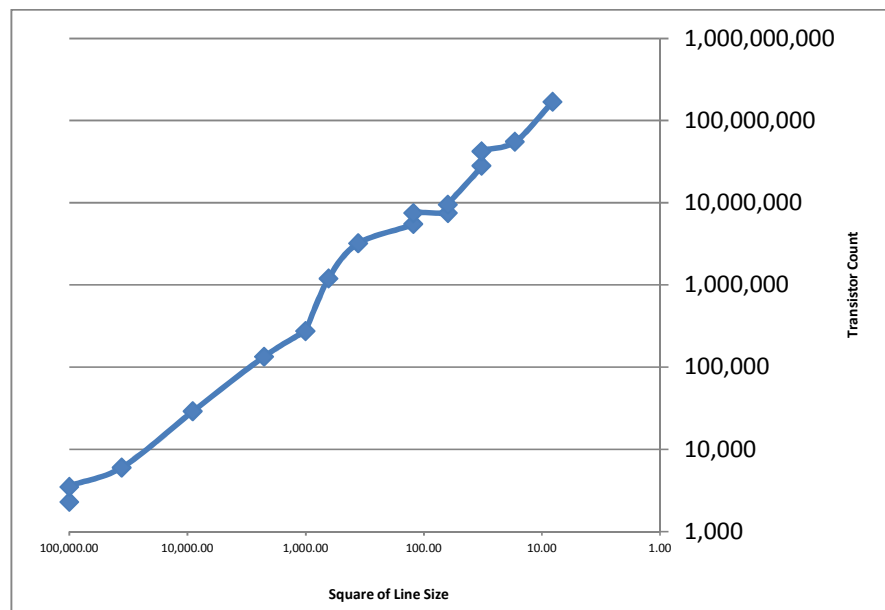
Here is the obligatory chart, illustrating Moore's law. The transistor count has grown by a factor of about 73,500 in about 33.25 years; doubling about every 2 years.



Two other charts are equally interesting. The first shows the increase in clock speed as a function of time. Note that this chart closely resembles the Moore's Law chart.



The second chart shows an entirely expected relationship between the line size on the die and the number of transistors placed on it. The count appears to be a direct function of the size.



Dual Core and Multi-Core CPU's

Almost all of the VLSI central processors on a chip that we have discussed to this point could be called “single core” processors; the chip had a single CPU along with some associated cache memory. Recently there has been a movement to multi-core CPU's, which represent the placement of more than one CPU on a single chip. One might wonder why this arrangement seems to be preferable to a single faster CPU on the chip.

In order to understand this phenomenon, we quote from a 2005 advertisement by Advanced Micro Devices (AMD), Inc. for the AMD 64 Opteron chip, used in many blade servers. This article describes chips by the linear size of the thinnest wire trace on the chip; a 90nm chip has features with a width of 90 nanometers, or 0.090 microns. Note also that it mentions the issue of CPU power consumption; heat dissipation has become a major design issue.

“Until recently, chip suppliers emphasized the frequency of their chips (“megahertz”) when discussing speed. Suddenly they’re saying megahertz doesn’t matter and multiple cores are the path to higher performance. What changed? Three factors have combined to make dual-core approaches more attractive.”

“First, the shift to 90nm process technology makes dual-core possible. Using 130nm technology, single-core processors measured about 200 mm², a reasonable size for a chip to be manufactured in high-volumes. A dual-core 130nm chip would have been about 400 mm², much too large to be manufactured economically. 90nm technology shrinks the size of a dual-core chip to under 200 mm² and brings it into the realm of possibility.”

“Second, the shift to dual-core provides a huge increase in CPU performance. Increasing frequency by 10 percent (often called a “speed bump”) results in at best a 10 percent boost in performance; two speed bumps can yield a 20 percent boost. Adding a second core can boost performance by a factor of 100 percent for some workloads.”

“Third, dual-core processors can deliver far more performance per watt of power consumed than single core designs, and power has become a big constraint on system design. All other things being equal, CPU power consumption increases in direct relation to clock frequency; boosting frequency by 20 percent boosts power by a least 20 percent. In practice, the power consumption situation is even worse. Chip designers need extra power to speed up transistor performance in order to attain increased clock frequencies; a 20 percent boost in frequency might require a 40 percent boost in power.”

“Dual-core designs apply this principle in reverse; they lower the frequency of each core by 20 percent, so that both cores combined use only a tad more power than a single-core at the higher frequency. This means the dual-core processors showing up this spring [2006?] will actually boost performance by a factor of approximately 1.7 over single core designs that fit in the same thermal envelope.

You can read more about dual-core technology at: **www.amd.com/dual-core**.”

Dual-core and quad-core systems are becoming quite popular. While there seems to be no theoretical upper limit on the number of cores (CPU's) on a chip, we may expect the practical limit to be in the range 8 to 16.

The State of the Art in 2008

The current state of the art can be illustrated by examination of a number of web sites for companies that sell computers. We consider the following examples.

The PowerEdge™ 2900 by Dell can be configured with 2 Quad-Core Intel Xenon processors, each operating at 3.16 GHz. The memory can range from 1 GB to 48 GB of 667 MHz synchronous DRAM, with error correcting code capability. The standard hard drives are either 3.5" SAS (operating at 15,000 RPM) with a maximum capacity of 450 GB per drive or a 3.5" SATA (operating at 7,200 RPM) with a maximum capacity of 1 TB per drive. The maximum internal disk storage is either 8 TB for SATA or 10 GB for SAS.

The latest IBM zSeries Enterprise server, the z10, is built on a quad-core processor operating at 4.46 GHz. Available systems can support up to 64 processors and from 16 GB to 1.5 TB of primary memory. Standard system storage options include a 9-track magnetic tape and a number of high-performance disk systems, such as the DS8300 which can store up to 512 TB, using 1024 500-GB FATA drives and 256 GB of buffer memory.

Another development of note is the **blade server**, which is a computer especially packaged to be rack mounted in a **blade enclosure**. There are many manufacturers of these devices. Blade server systems first appeared in 2001, with the introduction of the System 324 by RLX, Inc. This system could mount 324 blade servers in one enclosure about 74 inches high. Early blade servers delivered very interesting performance, but at a considerable downside. A typical rack of servers would consume about 30 kilowatts of power, necessarily radiating that as 30 kilowatts of heat that needed a good HVAC system to remove.

Modern blade servers are designed for lower power consumption. In the figure below, we see an IBM blade server on the right and a blade enclosure by Sun Microsystems, Inc. on the left. Sun describes the blade enclosure as follows.

“What makes IT organizations choose Sun is the additional power and cooling efficiencies they get with the Sun Blade 6000 Modular System over similar products from competitors. The Sun Blade 6000 Modular System is designed with strict front-to-back cooling, straight airflow, intelligent fan speed control, and better algorithms for maintaining sufficient airflow in compromised situations such as operating with failed fans. More efficient cooling means fewer watts spent on the cooling subsystem, a benefit that is amplified by lower CPU power consumption resulting from lower operating temperatures.”



**Sun Blade 6000
Modular System**



IBM HS20 Blade Server

Figure: Blade Enclosure and a Computer Configured as a Blade Server

Supercomputers: Their Rise, Fall and Reemergence

High-performance computing has always been of great importance in a wide variety of scientific and other disciplines. The emphasis on high performance gave rise to the supercomputer. While there have been many attempts to give the term “supercomputer” a precise definition, the common practice is just to equate the term with the list of computers designed and developed by Seymour Cray, first of the Control Data Corporation and then of the Cray Computer Corporation (a company he founded).

The first supercomputer is universally acknowledged to be the CDC-6600, introduced in 1964; the first unit was delivered to the National Center for Atmospheric Research in December 1965. The design was very simple, with every feature focused on one application. Unlike the IBM System/360 (introduced about the same time), the CDC-6600 made no pretense of being a general purpose computer; it was just a number crunching workhorse.

Much to the annoyance of the president of IBM, the CDC-6600 was significantly faster than the IBM 7094, up to then the standard machine for scientific computing.

The CDC 6600 is often called the first RISC (Reduced Instruction Set Computer), due to the simplicity of its instruction set. The reason for its simplicity was the desire for speed. Cray also put a lot of effort into matching the memory and I/O speed with the CPU speed. As he later noted, “Anyone can build a fast CPU. The trick is to build a fast system.”

The CDC 6600 led to the more successful CDC 7600. The CDC 8600 was to be a follow-on to the CDC 7600. While an excellent design, it proved too complex to manufacture successfully, and was abandoned. This feature will be seen again in future designs by Mr. Cray; he simply wanted the fastest machine that could be developed, almost without regard to the practicalities associated with its manufacture.

Cray left Control Data Corporation in 1972 to found Cray Research, based in Chippewa Falls, Wisconsin. The first Cray-1 was installed at the Los Alamos National Laboratory in 1976. According to the official Cray history web site [R54],

“It boasted a world-record speed of 160 million floating-point operations per second (160 megaflops) and an 8 megabyte (1 million word) main memory. ... In order to increase the speed of this system, the Cray-1 had a unique “C” shape which enabled integrated circuits to be closer together. No wire in the system was more than four feet long. To handle the intense heat generated by the computer, Cray developed an innovative refrigeration system using Freon.”

In addition to being the fastest standard computer of its day, the Cray-1 offered a very fast vector processor. The standard processor, by comparison, was called a scalar processor. Most scientific codes of the day made very heavy use of vectors, which are really just groupings of scalars (real numbers, integers, etc.). Consider this FORTRAN code.

```
DO 100 J = 1, 10
100  A[J] = B[J] + C[J]
```

This is just a vector addition; the elements of a 10-element vector B being added to those of a 10-element vector C and deposited in a 10-element vector A. Using a standard scalar processor, this would require ten distinct additions. In a vector processor, it essentially required only one addition. This proved a great advance in processing speed.

At over \$8 million a copy, the Cray-1 was called the “world’s most expensive love seat”. Note its appearance in the figure below, the low “seats” are really power supplies.



The Cray-1 at the Deutsches Museum in Munich Germany

The fundamental tension at Cray Research, Inc. was between Seymour Cray’s desire to develop new and more powerful computers and the need to keep the cash flow going.

Seymour Cray realized the need for a cash flow at the start. As a result, he decided not to pursue his ideas based on the CDC 8600 design and chose to develop a less aggressive machine. The result was the Cray-1, which was still a remarkable machine.

With its cash flow insured, the company then organized its efforts into two lines of work.

1. Research and development on the CDC 8600 follow-on, to be called the Cray-2.
2. Production of a line of computers that were derivatives of the Cray-1 with improved technologies. These were called the X-MP, Y-MP, etc.

The X-MP was introduced in 1982. It was a dual-processor computer with a 9.5 nanosecond (105 MHz) clock and 16 to 128 megawords of static RAM main memory. A four-processor model was introduced in 1984 with a 8.5 nanosecond (118 MHz) clock.

The Y-MP was introduced in 1988, with up to eight processors that used VLSI chips. It had a 32-bit address space, with up to 64 megawords of static RAM main memory. It was the first commercial computer to exceed 1 GigaFlop (10^9 floating point operations per second), achieving a sustained speed of 2.3 GigaFlops [R54].

While his assistant, Steve Chen, oversaw the production of the commercially successful X-MP and Y-MP series, Seymour Cray pursued his development of the Cray-2, a design based on the CDC 8600, which Cray had started while at the Control Data Corporation. The original intent was to build the VLSI chips from gallium arsenide (GaAs), which would allow much faster circuitry. The technology for manufacturing GaAs chips was not then mature enough to be useful as circuit elements in a large computer.

The Cray-2 was a four-processor computer that had 64 to 512 megawords of 128-way interleaved DRAM memory. The computer was built very small in order to be very fast, as a result the circuit boards were built as very compact stacked cards.



Due to the card density, it was not possible to use air cooling. The entire system was immersed in a tank of Fluorinert™, an inert liquid intended to be a blood substitute. When introduced in 1985, the Cray-2 was not significantly faster than the Y-MP. It sold only thirty copies, all to customers needing its large main memory capacity.

In 1989, Cray left the company in order to found Cray Computers, Inc. His reason for leaving was that he wanted to spend more time on research, rather than just churning out the very profitable computers that his previous company was manufacturing. This led to an interesting name game:

Cray Research, Inc.	producing a large number of commercial computers
Cray Computer, Inc.	mostly invested in research on future machines.

The Cray-3 was the first project Mr. Cray envisioned for his new company. This was to be a very small computer that fit into a cube one foot on a side. Such a design would require retention of the Fluorinert cooling system. It would also be very difficult to manufacture as it would require robotic assembly and precision welding. It would also have been very difficult to test, as there was no direct access to the internal parts of the machine.

The Cray-3 had a 2 nanosecond cycle time (with a 500 MHz clock). It used Gallium Arsenide semiconductor circuits (with a 80 picosecond switching time), which were five times faster than the silicon semiconductors used in the Cray-2. A single processor machine would have a performance of 948 megaflops; the 16-processor model would have operated at 15.2 gigaflops. A single-processor Cray-3 was delivered in 1993, but the 16-processor model was never delivered. The Cray-4, a smaller version of the Cray-3 with a 1 GHz clock was in development when Cray Computer, Inc. went bankrupt in 1995. Seymour Cray died on October 5, 1996 after an automobile accident.

In the end, the development of traditional supercomputers ran into several problems.

1. The end of the cold war reduced the pressing need for massive computing facilities. Some authors have linked the “computer technology race” to the nuclear arms race.
2. The rise of microprocessor technology allowing much faster and cheaper processors.
3. The rise of VLSI technology, making multiple processor systems more feasible.

Massively Parallel Processors (MPP)

In the late 1980's and 1990's, the key question in high-performance computer development focused on a choice: use a few (2 to 16) very high performance processors or a very large number of modified VLSI standard processors. As Seymour Cray put it: "If you were plowing a field, which would you rather use: Two strong oxen or 1024 chickens".

Although Seymour Cray said it more colorfully, there were many objections to the transition from the traditional vector supercomputer (with a few processors) to the massively parallel computing systems that replaced it. The key issue in assessing the commercial viability of a multiple-processor system is the speedup factor; how much faster is a processor with N processors. Here are two opinions from the 1984 IEEE tutorial on supercomputers.

"The speedup factor of using an n -processor system over a uniprocessor system has been theoretically estimated to be within the range $(\log_2 n, n/\log_2 n)$. For example, the speedup range is less than 6.9 for $n = 16$. Most of today's commercial multiprocessors have only 2 to 4 processors in a system."

"By the late 1980s, we may expect systems of 8–16 processors. Unless the technology changes drastically, we will not anticipate massive multiprocessor systems until the 90s."

Two things facilitated the introduction of the Massively Parallel Processor systems.

1. The introduction of fast and cheap processor chips, such as the Pentium varieties.
2. The discovery that many very important problems could be solved by algorithms that displayed near "linear speedup". Thus, multiplying the number of processors by a factor of ten might well speed up the solution by a factor of nine.

Here is a picture of the Cray XT-5, one of the later and faster products from Cray, Inc. [R83] It is a MPP (Massively Parallel Processor) system, launched in November 2007.



This is built from a number of Quad-Core AMD Opteron™ processor cores, which are variants of the Intel Pentium chip. At least one model of this computer, as well as a MPP model from IBM, have sustained a speed of 1 Teraflop (10^{12} floating point operations per second) on real problems of significance to research scientists. The supercomputer is back.