Chapter 16: Direct Conversions Between EBCDIC and Fullword Formats

This chapter presents a discussion of direct conversions between digits in the EBCDIC format and binary integers stored in the 32–bit two’s–complement format. This material is presented within the context of an academic exercise focused on gaining a more complete understanding of the basic principles involved. In reality, a program is much more likely to use the existing tools (PACK, CVB, CVD, and ED) provided by the S/370 assembler.

In other words, the goal of this chapter is not to add to the student’s “bag of assembler tricks” but to add to the student’s knowledge.

Two’s–Complement Binary Format
Binary integer data are stored on the System/370 in two basic formats.

1. **Halfword** 16–bit two’s–complement integers
2. **Fullword** 32–bit two’s–complement integers.

The **halfword format** is conventionally represented by four hexadecimal digits, which occupy two bytes of storage. A properly aligned halfword has an address that is a multiple of 2. The range of values that can be represented is from \(-32,768\) to 32,767 inclusive. The print representation of a halfword integer contains at most five digits.

The **fullword format** is conventionally represented by eight hexadecimal digits, which occupy four bytes of storage. A properly aligned halfword has an address that is a multiple of 4. The range of values that can be represented is from \(-2,147,483,648\) to 2,147,483,647. The print representation of a fullword integer contains at most ten digits.

The name “two’s–complement” refers to the manner of storing negative integers. The student should review the material in chapter 4 of this textbook, especially that on conversion from decimal to binary format, binary to decimal format, and taking the two’s–complement. Here is a very short presentation on the topic.

The positive decimal number 165 can be represented in hexadecimal as \(X'\text{A5}'\). As an eight bit binary number, this is \(1010 0101\). We now consider the representation of the negative decimal number \(-165\). In order to give the binary representation, we must specify the format.

As a 16–bit number \(+165\) is \(0000 0000 1010 0101\) or \(X'\text{00A5}'\)
take the one’s complement \(1111 1111 0101 1010\)
add one to get the result \(1111 1111 0101 1011\) or \(X'\text{FF5A}'\)
This last number is the binary representation of \(-165\) as a 16–bit integer.

As a 32–bit number \(+165\) is \(0000 0000 0000 0000 0000 1010 0101\)
take the one’s complement \(1111 1111 1111 1111 1111 1111 0101 1010\)
add one to get the result \(1111 1111 1111 1111 1111 1111 0101 1011\)
This last number, also represented as \(X'\text{FFFF FF5A}'\), is the binary representation of \(-165\) as a 32–bit binary integer.

Integers are converted from fullword (32 bits or 4 bytes) to halfword (16 bits or 2 bytes) format by copying the rightmost two bytes, represented by four hexadecimal digits. If the number is too large in magnitude for the halfword format, it is truncated.
Integers are converted from halfword (16 bit or 2 bytes) to fullword (32 bits or 4 bytes) format by sign extension. This process insures that the sign of the number is preserved.

+165 in 16 bits is \(0000 \ 0000 \ 1010 \ 0101\)
+165 in 32 bits is \(0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 1010 \ 0101\)

–165 in 16 bits is \(1111 \ 1111 \ 0101 \ 1011\)
–165 in 32 bits is \(1111 \ 1111 \ 1111 \ 1111 \ 1111 \ 1111 \ 0101 \ 1011\)

In each case, it is the leftmost bit in the 16–bit (halfword) representation that is copied to the leftmost 16 bits added when moving to the 32–bit (fullword) format.

The assumption is that the binary number to be considered will be stored in a general–purpose register, such as R7. The register might be loaded by an instruction such as one of the two following.

\[
\begin{align*}
\text{L } & \text{R7,FW1} & \text{Load R7 from a fullword in memory} \\
\text{LH } & \text{R7,HW1} & \text{Load R7 from a halfword in memory, and sign extend to a fullword format.}
\end{align*}
\]

The goal of the input program will be to convert from the EBCDIC digit representation, which is really just a sequence of character codes, into a binary number in a register.

**EBCDIC Representation of Digits**

When digits are read in from an input device, they are treated as character data that only incidentally have numeric value. These must be converted to a numeric format.

The EBCDIC codes of interest in the representation of integer data are the following.

<table>
<thead>
<tr>
<th>Code</th>
<th>Digit</th>
</tr>
</thead>
<tbody>
<tr>
<td>X<code>F0</code></td>
<td>0</td>
</tr>
<tr>
<td>X<code>F1</code></td>
<td>1</td>
</tr>
<tr>
<td>X<code>F2</code></td>
<td>2</td>
</tr>
<tr>
<td>X<code>F3</code></td>
<td>3</td>
</tr>
<tr>
<td>X<code>F4</code></td>
<td>4</td>
</tr>
<tr>
<td>X<code>F5</code></td>
<td>5</td>
</tr>
<tr>
<td>X<code>F6</code></td>
<td>6</td>
</tr>
<tr>
<td>X<code>F7</code></td>
<td>7</td>
</tr>
<tr>
<td>X<code>F8</code></td>
<td>8</td>
</tr>
<tr>
<td>X<code>F9</code></td>
<td>9</td>
</tr>
</tbody>
</table>

The two other codes of interest are X`40` for the space and X`60` for the minus sign.

**Print Representation of Integers**

It goes without saying that the print representation of any integer will involve the use of EBCDIC characters, especially the ones listed just above. What must be considered is how to present negative integers. Consider the negative integer 165 to be printed as four digits.

The standard algebraic way to do this is \(-165\).

A less used way is to print it in this form \(-165\).

A way commonly seen in mainframe programs is as follows \(165-\).

The last way, though appearing strange, is quite easy to program. For this reason, many assembler language programs will use the “postfix minus sign” for negative numbers. The second way involves a bit more code to produce, and the first way considerably more code. It is this algebraically correct representation that is our goal in this chapter.
NUMIN: A Program to Input Binary Integers
Numeric data are input into a computer in a three step process.

1. The data are read in as a sequence of characters.
   For the IBM System/360, the characters are encoded as EBCDIC.
2. The data are converted to the proper form for numeric use.
3. The data are stored, either in memory or general–purpose registers,
   for use in computations.

We shall focus on the input of integer data to be stored in one of the general–purpose
registers. As an arbitrary constraint, we shall limit the numbers to 9 digits, though the
numbers are allowed to be smaller.

Note that any possible nine–digit integer can be stored as a 32–bit fullword. While it is the
case that some ten–digit numbers can be stored as a fullword, this does not hold for all such
numbers; for example:
   The ten digit number 2, 100, 000, 000 can be converted to fullword format.
   The ten digit number 2, 200, 000, 000 cannot be converted to fullword format.

It is for this reason that our code will focus on numbers with a maximum of nine digits,
represented by ten characters, allowing for an optional sign character.

NUMIN: The Scenario
Remember that input should be viewed as a card image of 80 columns. Consider a field of
N characters found beginning in column M.

Suppose that the leftmost byte in this array is associated with the label \texttt{CARDIN}. The
leftmost byte in the range of interest will be denoted by the label \texttt{CARDIN+M}. Elements in
this range will be referenced using an index register as \texttt{CARDIN+M(Reg)}, where the
number in parentheses represents the index register to be used.

Our specific example will assume the following:

1. The character field to hold the integer occupies ten columns on the card,
   beginning in column 20 and running through column 29.
2. The number is right justified. If negative, the number has a leading minus sign.
3. An entirely blank field is accepted as representing the number zero.
NUMIN: The Standard Approach
We begin this set of notes by recalling a more standard approach to conversion from a sequence of EBCDIC characters to a binary number in a register. This sample code will assume that all numbers are non-negative.

Here are some data declarations that are used in the code. Note that the data declaration seems to call for ten digits. Here the assumption will be that the input has at least one leading space and at most nine numeric digits with no sign.

* THE CHARACTERS FOR INPUT ARE FOUND BEGINNING
* AT CARDIN+20 THROUGH CARDIN+29. NO MINUS SIGN.

DIGITSIN DS CL10 TEN BYTES TO HOLD 10 CHARACTERS
PACKEDIN DS PL6 SIX BYTES HOLD 11 DIGITS
PACKDBL DS D DOUBLE WORD TO HOLD PACKED

Here is the code that uses the above data structures.

MVC DIGITSIN(10),CARDIN+20 GET 10 CHARACTERS
PACK PACKEDIN,DIGITSIN CONVERT TO PACKED
ZAP PACKDBL,PACKEDIN FORMAT FOR CVB
CVB R7,PACKDBL BINARY INTO R7.

NUMIN: The Strategy
The figure below shows the part of the 80–column card image that contains the digits to be interpreted. We now discuss the strategy to be followed in our direct conversion routine.

The algorithm works as follows:
1. It initializes an output register to 0. Arbitrarily, I choose R7.
2. It scans left to right, looking for a nonblank character.

Assuming that a nonblank character is found in this field, it does the following.

3. If the character is a minus sign, set a flag that the number is negative and continue the scan.
4. If the number is a digit, process it. If not a digit or “–”, ignore it.

One problem of this code is typical of most sample code. In an attempt to focus on one point, the code ignores all error processing. Just be aware of the fact.
NUMIN: EXAMPLE

Consider processing the number represented by the digit string “9413”. We shall illustrate the process used by our conversion routine.

In this example, let

- \( N \) be the value of the number,
- \( D \) be the digit read in, and
- \( V \) be the numeric value of that digit.

Start with \( N = 0 \).

Read in \( D = \text{"9"} \). Convert to \( V = 9 \). \( N = N \times 10 + V = 0 \times 10 + 9 = 9 \)

Read in \( D = \text{"4"} \). Convert to \( V = 4 \). \( N = N \times 10 + V = 9 \times 10 + 4 = 94 \)

Read in \( D = \text{"1"} \). Convert to \( V = 1 \). \( N = N \times 10 + V = 94 \times 10 + 1 = 941 \)

Read in \( D = \text{"3"} \). Convert to \( V = 3 \). \( N = N \times 10 + V = 941 \times 10 + 3 = 9413 \)

The integer value of this string is 9413.

Review of the Instructions: LCR and IC

The code below will use two instructions that should be reviewed at this point. These are LCR (Load Complement Register) and IC (Insert Character).

**Load Complement Register: LCR R1, R2**

This loads register R1 with the negative (two’s–complement) of the value in register R2. This is a convenient way to change the sign of the integer in a register; set the value in the register equal to the negative of the value now there.

**Insert Character: IC R8, CARDIN+20(R3) GET THE DIGIT**

This inserts the eight bits of the EBCDIC character into the low order 8 bits (bits 24 – 31) of the destination register. The other bits are not changed.

There are many interesting uses of this instruction. I elect to use this to set the value in the register equal to the value of a digit. Thus if the character with EBCDIC representation X‘F7’ is in storage, I can set the value in the register to 7.
Placing the Numerical Value of a Digit in a Register
The first thing to do is get the EBCDIC code into the register. My solution uses the IC (Insert Character) instruction.

```
SR  R8,R8    CLEAR R8
IC  R8,CARDIN+20(R3) GET THE DIGIT
S   R8,=X'F0' CONVERT TO VALUE OF DIGIT
```

In order to be sure that register R8 contains the EBCDIC code for the digit, I first clear the register to zero and then move the character. This step guarantees that bits 0–23 of the register are 0 and that the value in the register, taken as a 32–bit fullword, is the EBCDIC code for the digit. I then subtract the value of the EBCDIC code for ‘0’ to get the value.

Another way to do this is load the register and use the logical instruction, with mnemonic N, to mask out all but the last hexadecimal digit. Here is the code.

```
IC  R8,CARDIN+20(R3) GET THE DIGIT
N   R8,=X'F'          
```

I now present my algorithm in fragments of code. We start with the beginning code. Each fragment will be listed along with its associated data declarations. This first code fragment just clears the result registers and checks to see if the input field, in the ten columns beginning at CARDIN+20, is all blanks.

If it is all blanks, the routine interprets the field as containing a 0 and returns.

```
NUMIN  SR  R7,R7      SET R7, THE RESULT, TO 0
       SR  R6,R6      CLEAR HIGH-ORDER PRODUCT
       MVI  THESIGN,C 'P' DEFAULT TO POSITIVE
       CLC  CARDIN+20(10),SPACE10 IS THE INPUT ALL BLANKS
       BE DONE IF SO, JUST EXIT WITH

* MORE CODE HERE

* 0123456789 BE SURE OF THE COUNT BELOW
SPACE10 DC CL ' ' JUST TEN SPACES
THESIGN DS CL1
```

The next part scans left to right looking for a non–blank character, which should be there. If none is found, it just quits. Admittedly, this should not happen, as we have tested and found at least one non–blank character in the input. This is defensive coding.

```
* NOW SCAN LEFT TO RIGHT TO FIND FIRST NON-BLANK.
* USE BXLE WITH REGISTER PAIR (R4,R5).

* SR  R3,R3 CLEAR INDEX USED TO SCAN THE INPUT CHARACTER ARRAY

* LA  R4,1 SET INCREMENT TO 1
LA  R5,9 OFFSET 9 IS THE LAST DIGIT

SCAN1 CLI  CARDIN+20(R3),C ' ' DO WE HAVE A SPACE?
BNE  NOTBLANK NO, IT MAY BE A DIGIT
BXLE R3,R4,SCAN1 ITS BLANK, LOOK AT NEXT
B    DONE ALL BLANKS, WE ARE DONE
```
This next section of code checks the first non-blank character. If it is a minus sign, the sets a flag, which would be a Boolean in a high-level language. Here it is just the character “N”.

If the first non-blank character is a minus sign, then the next character is assumed to be the first digit. The index value is incremented by 1 to address the character after the “-”.

If the first non-blank character is not a minus sign, it is assumed to be a digit and processed as one. Note however that the processing loop explicitly makes two tests and processes the character only if it is not less than ‘0’ and not greater than “9”.

```
* AT THIS POINT, R3 IS THE INDEX OF THE NON-BLANK
* CHARACTER. THE VALUES IN (R4,R5) ARE STILL VALID.
* IN PARTICULAR R4 STILL HAS VALUE 1.
*
NOTBLANK CLI CARDIN+20(R3),C'-' DO WE HAVE A MINUS SIGN?
BNE ISDIG
MVI THESIGN,C 'N' NOTE THE SIGN AS NEGATIVE
AR R3,R4 ADD 1 TO VALUE IN R3.
CR R3,R5 R3 HAS BEEN INCREMENTED
BH DONE QUIT IF IT IS TOO BIG.
```

At this point, we know that `CARDIN+20(R3)` references a non-blank character that is in the range of card columns that might contain a digit. Here is the conversion loop. Note that the first four lines check to see if the character is a digit by performing two tests equivalent to the compound inequality ‘0’ \( \leq \) Code \( \leq \) ‘9’. If the character is not a digit, it is ignored and a branch to the end of the loop is taken.

```
ISDIG CLI CARDIN+20(R3),C '0' IS IT A DIGIT
BL LOOP NO - CODE < '0'
CLI CARDIN+20(R3),C '9' AGAIN, IS IT A DIGIT?
BH LOOP NO - CODE > '9'
M R6,=F'10' MULTIPLY (R6,R7) BY 10
SR R8,R8 CLEAR R8
IC R8,CARDIN+20(R3) GET THE DIGIT
S R8,=X'F0' CONVERT TO VALUE OF DIGIT
AR R7,R8 ADD TO THE PRODUCT
LOOP BXLE R3,R4,ISDIG END OF THE LOOP
CLI THESIGN,C 'N' WAS THE INPUT NEGATIVE
BNE DONE IT IS NOT NEGATIVE
LCR R7,R7 TAKE 2'S COMPLEMENT
DONE * HERE R7 CONTAINS THE BINARY VALUE
```
Here is the complete code for NUMIN.

```assembly
NUMIN      SR R7, R7           SET R7, THE RESULT, TO 0
SR R6, R6          CLEAR HIGH-ORDER PRODUCT
MVI THESIGN, C'P'      DEFAULT TO POSITIVE
CLC CARDIN+20(10), SPACE10  IS THE INPUT ALL BLANKS
BE DONE            IF SO, JUST EXIT WITH
  THE VALUE SET TO 0.

  *  
  *  
  *  NOW SCAN LEFT TO RIGHT TO FIND FIRST NON-BLANK.
  *  USE BXLE WITH REGISTER PAIR (R4, R5).
  *  
  SR R3, R3          CLEAR INDEX USED TO SCAN
  *
LA R4, 1            THE INPUT CHARACTER ARRAY
LA R5, 9          SET INCREMENT TO 1
OFFSET 9 IS THE LAST DIGIT
SCAN1 CLI CARDIN+20(R3), C' '  DO WE HAVE A SPACE?
BNE NOTBLANK        NO, IT MAY BE A DIGIT
BXLE R3, R4, SCAN1   ITS BLANK. LOOK AT NEXT
B DONE              ALL BLANKS, WE ARE DONE

  *  
  *  AT THIS POINT, R3 CONTAINS THE INDEX OF THE NON-BLANK
  *  CHARACTER. THE VALUES IN (R4, R5) ARE STILL VALID.
  *  IN PARTICULAR R4 STILL HAS VALUE 1.
  *  
NOTBLANK CLI CARDIN+20(R3), C'-'  DO WE HAVE A MINUS SIGN?
  BNE ISDIG             NO
  MVI THESIGN, C'N'      NOTE THE SIGN AS NEGATIVE
  AR R3, R4              ADD 1 TO VALUE IN R3.
  CR R3, R5              R3 HAS BEEN INCREMENTED
  BH DONE              QUIT IF IT IS TOO BIG.
  *
ISDIG CLI CARDIN+20(R3), C'0'  IS IT A DIGIT
  BL LOOP              NO - CODE < '0'
  CLI CARDIN+20(R3), C'9'  AGAIN, IS IT A DIGIT?
  BH LOOP              NO - CODE > '9'
  M R6, =F'10'         MULTIPLY (R6, R7) BY 10
  SR R8, R8         CLEAR R8
  IC R8, CARDIN+20(R3)  GET THE DIGIT
  S R8, =X'F0'       CONVERT TO VALUE OF DIGIT
  AR R7, R8         ADD TO THE PRODUCT
LOOP BXLE R3, R4, ISDIG          END OF THE LOOP
  CLI THESIGN, C'N'    WAS THE INPUT NEGATIVE
  BNE DONE            IT IS NOT NEGATIVE
  LCR R7, R7       TAKE 2'S COMPLEMENT

  *  
DONE        * HERE R7 CONTAINS THE BINARY VALUE
  *
  0123456789  BE SURE OF THE COUNT BELOW
SPACE10    DC CL' '          JUST TEN SPACES
  THESIGN   DS CL1
```
Printing Packed Data

The standard solution to convert binary integer data into printable form uses two of the standard System/370 assembler language instructions.

- **CVD**: Converts the binary to packed decimal.
- **UNPK**: Converts the packed decimal to zoned decimal format.

The unpack command, UNPK, has an unfortunate side effect. Consider the decimal number 42, represented in binary in register R4.

**CVD R4,PACKOUT** produces the value in standard packed decimal format: **042C**.

This should be unpacked to the EBCDIC **F0 F4 F2**

**UNPK** produces the zoned format **F0 F4 C2**.

This prints as “**04B**”, because **0xC2** is the EBCDIC code for the letter ‘B’.

Here is the code that works.

```assembly
NUMOUT CVD R4,PACKOUT
UNPK THESUM,PACKOUT
MVZ THESUM+7(1),=X’F0’
* BR 8
PACKOUT DS PL8
```

**THESUM** has eight characters stored as eight bytes. The addresses are:

<table>
<thead>
<tr>
<th>SUM</th>
<th>SUM +1</th>
<th>SUM +2</th>
<th>SUM +3</th>
<th>SUM +4</th>
<th>SUM +5</th>
<th>SUM +6</th>
<th>SUM +7</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Hundreds</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>Tens</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Units</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Again, the expression **THESUM+7** is an address, not a value.

If **THESUM** holds **C’01234567’**, then **THESUM+7** holds **C’7’**.

A Problem with the Above Routine

Consider the decimal number –42, stored in a register in binary two’s–complement form.

- **CVD** produces **042D**
- **UNPK** produces **F0 F4 D2**

The above **MVZ** will convert this to **F0 F4 F2**, a positive number. There are some easy fixes that are guaranteed to produce the correct representation for a negative number.

Most of the fixes using CVD and UNPK depend on placing the minus sign to the right of the digits. So that the negative integer –1234 would be printed as “**1234–**”.

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Printed by Edward L. Bosworth, Ph.D.
My Version of NUMOUT (Number Out)
This routine avoids packed decimal numbers. We are given a binary number (negative or non–negative) in register R4.

1. Is the number negative?
   If so, set the sign to ‘–’ and take the absolute value.
   Otherwise, leave the sign as either ‘+’ or ‘ ’ (a blank).

We now have a non–negative number. Assume it is not zero.

2. Divide the number by 10, get a quotient and a remainder.
   The remainder will become the character output.

3. The remainder is a positive number in the range [0, 9].
   Add =X’F0’ to produce the EBCDIC code.

4. Place this digit code in the proper output slot.
Is the quotient equal to 0? If so, quit.
If it is not zero, place the quotient in the dividend and return to 2.

Here is a paper example of the proper execution of the algorithm. Consider the positive integer 9413. Do repeated division by 10 and watch the remainders.

9413 divided by 10:  Quotient = 941  Remainder = 3.  Generate digit “3”.
941 divided by 10:   Quotient = 94  Remainder = 1.  Generate digit “1”.
94 divided by 10:    Quotient = 9   Remainder = 4.  Generate digit “4”.
9 divided by 10:     Quotient = 0   Remainder = 9.  Generate digit “9”.

Quotient is zero, so the process stops.

As they are generated, the digits are placed right to left, so that the result will print as the string “9413”. We now investigate the specifications for the code.

NUMOUT: Specifications
The code processes a 32–bit two’s–complement integer, stored as a fullword in register R5 and prints it out as a sequence of EBCDIC characters. The specification calls for printing out at most 10 digits, each as an EBCDIC character. The sign will be placed in the normal spot, just before the number. For no particular reason, positive numbers will be prefixed with a “+”. I just thought I would do something different.

This will use repeated division, using the even–odd register pair (R4, R5), which contains a 64–bit dividend. As a part of our processing we shall insure that the dividend is a 32–bit positive number. In that case, the “high order” 32 bits of the number are all 0.

For that reason, we initialize the “high order” register, R4, to 0 and initialize the “low order” register, R5, to the absolute value of the integer to be output.

The EBCDIC characters output will be placed in a 12–byte area associated with the label CHARSOUT, at byte addresses CHARSOUT through CHARSOUT+11.
Review of the Instructions: LCR and STC

Load Complement Register: \texttt{LCR R1, R2}

This loads register R1 with the negative (two’s-complement) of the value in register R2. This is also used in my routine NUMIN.

Store Character: \texttt{STC R8,CHARSOUT(R3) PLACE THE DIGIT}

This transfers the EBCDIC character, with code in the low order 8 bits of the source register, to the target address. None of the bits in the register are changed.

The idea behind NUMOUT is to compute the numerical value of a digit in a source register, convert it to an EBCDIC code, and move it to the print line. The first part checks the sign of the integer in register R4 and sets the sign character appropriately.

Note that the first thing to do is clear the output field to that expected for a zero result.

\begin{verbatim}
NUMOUT MVI THESIGN,C’+’ DEFAULT TO 0
MVI THESIGN,C’-’ DEFAULT TO A PLUS SIGN
C R5,=F’0’ COMPARE R5 TO 0
BE DONE VALUE IS 0, NOTHING TO DO
BH ISPOS VALUE IS POSITIVE
MVI THESIGN,C’-’ PLACE A MINUS SIGN
LCR R5,R5 2’S COMPLEMENT R5 TO MAKE POS
ISPOS SR R4,R4 CLEAR REGISTER 4
\end{verbatim}

Here are some data declarations used with this part of the code.

\begin{verbatim}
* 123456789012
ZEROOUT DC C’ 0’ 11 SPACES AND A ZERO
CHARSOUT DS CL12 UP TO 11 DIGITS AND A SIGN
\end{verbatim}

Division (Specifically D – Divide Fullword)

This instruction divides a 64–bit dividend, stored in an even–odd register pair, by a fullword, and places the quotient and remainder back into the register pair.

This will use the even–odd register pair (R4, R5). The specifics of the divide instruction are as follows.

<table>
<thead>
<tr>
<th>Before division</th>
<th>R4</th>
<th>R5</th>
</tr>
</thead>
<tbody>
<tr>
<td>Dividend (high order 32 bits)</td>
<td>Dividend (low order 32 bits)</td>
<td></td>
</tr>
<tr>
<td>After division</td>
<td>Remainder</td>
<td>Quotient</td>
</tr>
</tbody>
</table>

There are specific methods to handle dividends that might be negative.

As we are considering only positive dividends, we ignore these general methods.
Our Example of Division
Start with a binary number in register R5.

We assume that register R4 has been cleared to 0, as this example is limited to a 32-bit positive integer. This code will later be modified to process the remainder, and store the result as a printable EBCDIC character.

Here is the broad outline of the conversion loop, called DIVIDE because it achieves the result by repeated division by ten.

```
DIVIDE  D  R4,=F'10'
DIVIDE (R4,R5) BY TEN

*   THE REMAINDER, IN R4, MUST BE PROCESSED AND STORED
*
SR  R4,R4
C  R5,=F'0'
BH  DIVIDE
```

Placing the Digits
At this point, our register and storage usage is as follows:
1. Register R3 will be used as an index register.
2. Register pair (R4, R5) is being used for the division.
3. Register pair (R6, R7) is reserved for use by the BXH instruction.

**CHARSOUT DS CL12** contains the twelve characters that form the print representation of the integer. The number 12 is arbitrary; it could be 10.

The strategy calls for first placing a digit in the units slot (overwriting the ‘0’) and then moving left to place other digits. To allow for a sign, no digit is to be placed in slot 0, at address **CHARSOUT**. The idea will be to place the character into a byte specified by **CHARSOUT(R3)**. The register is initialized at 11 and decremented by 1 using the BXH instruction. What the code actually does is increment R3 by the negative value –1.

```
0
```

The Digit Placement Code
Here is a sketch of the digit placement code. It must be integrated into the larger DIVIDE loop in order to make sense. The register pair (R6, R7) is used for the BXH instruction.

R6 holds the increment value
R7 holds the limit value

```
L  R6,=F'−1'
SR  R7,R7
L  R3,=F'11'
A  R4,=X'F0'
STC  R4,CHARSOUT(R3)
BXH  R3,R6,DIVIDE
MVC  CHARSOUS(R3),THESIGN
```

```
SET INCREMENT TO −1
CLEAR R7. LIMIT VALUE IS 0.
SET INDEX TO 11 FOR LAST DIGIT.
ADD TO GET EBCDIC CODE
PLACE THE CHARACTER
GO BACK TO TOP OF LOOP
PLACE THE SIGN
```
The Complete Divide Loop
Here is the complete code for the divide loop. Note the branch out of the loop. The loop exits either when the quotient is 0 or when ten digits have been placed.

```
L R6,=F'1'          SET INCREMENT TO -1
SR R7,R7            CLEAR R7. LIMIT VALUE IS 0.
L R3,=F'11'          SET INDEX TO 11 FOR LAST
                      DIGIT AT CHAROUT+11.

DIVIDE D R4,=F'10'   DIVIDE (R4,R5) BY TEN AND
A R4,=X'F0'          ADD X 'F0', THE CODE FOR '0'
                      TO GET EBCDIC CODE FOR DIGIT
STC R4,CHARSOUT(R3)  PLACE THE CHARACTER
SR R4,R4            CLEAR R4 FOR ANOTHER LOOP
C R5,=F'0'          CHECK THE QUOTIENT
BNH PUTSIGN         EXIT LOOP IF QUOTIENT <= 0
BXH R3,R6,DIVIDE    GO BACK TO TOP OF LOOP

PUTSIGN MVC CHAROUT(R3),THESIGN    PLACE THE SIGN

Here is the complete code for NUMOUT.

*THE FIRST PART SETS THE DEFAULTS AND PREPARES FOR A 0 OUTPUT

NUMOUT MVC CHAROUT,ZEROOUT    DEFAULT TO 0
MVI THESIGN,C+'       DEFAULT TO A PLUS SIGN
C R5,=F'0'          COMPARE R5 TO 0
BE DONE          VALUE IS 0, NOTHING TO DO
BH ISPOS        VALUE IS POSITIVE
MVI THESIGN,C'-       PLACE A MINUS SIGN
LCR R5,R5          2'S COMPLEMENT R5 TO MAKE POS
ISPOS SR R4,R4       CLEAR REGISTER 4

L R6,=F'1'          SET INCREMENT TO -1
SR R7,R7            CLEAR R7. LIMIT VALUE IS 0.
L R3,=F'11'          SET INDEX TO 11 FOR LAST
                      DIGIT AT CHAROUT+11.

DIVIDE D R4,=F'10'   DIVIDE (R4,R5) BY TEN AND
A R4,=X'F0'          ADD X 'F0', THE CODE FOR '0'
                      TO GET EBCDIC CODE FOR DIGIT
STC R4,CHARSOUT(R3)  PLACE THE CHARACTER
SR R4,R4            CLEAR R4 FOR ANOTHER LOOP
C R5,=F'0'          CHECK THE QUOTIENT
BNH PUTSIGN         EXIT LOOP IF QUOTIENT <= 0
BXH R3,R6,DIVIDE    GO BACK TO TOP OF LOOP

PUTSIGN MVC CHAROUT(R3),THESIGN    PLACE THE SIGN IN THE SPOT
                      FOR STANDARD ALGEBRA

* CODE HERE FOR RETURN FROM SUBROUTINE
```